2-coherent and 2-convex Conditional Lower Previsions

Renato Pelessoni*1 and Paolo Vicig†1

1DEAMS “B. de Finetti”, University of Trieste, Piazzale Europa 1, I-34127 Trieste, Italy

June 20, 2016

Abstract

In this paper we explore relaxations of (Williams) coherent and convex conditional previsions that form the families of \( n \)-coherent and \( n \)-convex conditional previsions, at the varying of \( n \). We investigate which such previsions are the most general one may reasonably consider, suggesting (centered) 2-convex or, if positive homogeneity and conjugacy is needed, 2-coherent lower previsions. Basic properties of these previsions are studied. In particular, we prove that they satisfy the Generalized Bayes Rule and always have a 2-convex or, respectively, 2-coherent natural extension. The role of these extensions is analogous to that of the natural extension for coherent lower previsions. On the contrary, \( n \)-convex and \( n \)-coherent previsions with \( n \geq 3 \) either are convex or coherent themselves or have no extension of the same type on large enough sets. Among the uncertainty concepts that can be modelled by 2-convexity, we discuss generalizations of capacities and niveloids to a conditional framework and show that the well-known risk measure Value-at-Risk only guarantees to be centered 2-convex. In the final part, we determine the rationality requirements of 2-convexity and 2-coherence from a desirability perspective, emphasising how they weaken those of (Williams) coherence.

Keywords. Williams coherence, 2-coherent previsions, 2-convex previsions, Generalised Bayes Rule.

Acknowledgement

*NOTICE: This is the authors’ version of a work that was accepted for publication in the International Journal of Approximate Reasoning. Changes resulting from the publishing process, such as peer review, editing, corrections, structural...
formatting, and other quality control mechanisms may not be reflected in this document. Changes may have been made to this work since it was submitted for publication. A definitive version was subsequently published in the International Journal of Approximate Reasoning, vol. 77, October 2016, pages 66–86, doi:10.1016/j.ijar.2016.06.003 © Copyright Elsevier http://www.sciencedirect.com/science/article/pii/S0888613X16300792.

© 2016. This manuscript version is made available under the CC-BY-NC-ND 4.0 license http://creativecommons.org/licenses/by-nc-nd/4.0/

1 Introduction

In his influential book Statistical Reasoning with Imprecise Probabilities [21], P. Walley developed a behavioural approach to imprecise probabilities (and previsions) extending de Finetti’s [5] interpretation of coherent precise previsions. Operationally, this was achieved through a relaxation of de Finetti’s betting scheme. In fact, following de Finetti, $P$ is a coherent precise prevision on a set $S$ of gambles if and only if for all $m, n \in \mathbb{N}_0$, $s_1, \ldots, s_m, r_1, \ldots, r_n \geq 0$, $X_1, \ldots, X_m$, $Y_1, \ldots, Y_n \in S$, defining $G = \sum_{i=1}^{m} s_i(X_i - P(X_i)) - \sum_{j=1}^{n} r_j(Y_j - P(Y_j))$, it holds that $\sup G \geq 0$. The terms $s_i(X_i - P(X_i))$, $-r_j(Y_j - P(Y_j))$ are proportional (with coefficients or stakes $s_i$, $r_j$) to the gains arising from, respectively, buying $X_i$ at $P(X_i)$ or selling $Y_j$ at $P(Y_j)$. A coherent lower prevision $P$ on $S$ may be defined in a similar way, just restricting $n$ to belong to $\{0, 1\}$. This means that the betting scheme is modified to allow selling at most one gamble. Several other betting scheme variants have been investigated in the literature, either extending coherence for lower previsions (conditional lower previsions) or weakening it (previsions that are convex, or avoid sure loss). In particular, a convex lower prevision is defined introducing a convexity constraint $n = 1, \sum_{i=1}^{m} s_i = r_1 = 1$ in the betting scheme. In [21, Appendix B] $n$-coherent previsions are studied, as a different relaxation of coherence.

In this paper, we explore further variations of the behavioural approach /betting scheme: $n$-coherent and $n$-convex conditional lower previsions, formally defined later on as generalisations of the $n$-coherent (unconditional) previsions in [21]. Our major aims are:

a) to explore the flexibility of the behavioural approach and its capability to encompass different uncertainty models;

b) to point out which are the basic axioms/properties of coherence which hold even for much looser consistency concepts.

Referring to b) and with a view towards the utmost generality, we shall mainly concentrate on the extreme quantitative models that can be incorporated into a
(modified) behavioural approach. This does not imply that these models should be regarded as preferable to coherent lower previsions. On the contrary they will not, as far as certain questions are concerned. For instance, inferences will typically be rather vague. However, it is interesting and somehow surprising to detect that certain properties like the Generalised Bayes Rule must hold even for such models, or that they can be approached in terms of desirability.

\(N\)-coherence and \(n\)-convexity may be naturally seen as relaxations of, respectively, (Williams) coherence and convexity. These and other preliminary concepts are recalled in Section 2. Starting from the weakest reasonably sound consistency concepts, we explore 2-convex lower previsions in Section 3. In Section 3.1 we characterise them by means of axioms, on a special set of conditional gambles generalising a linear space and termed \(D_{LIN}\) (Definition 2, Proposition 2). Interestingly, it turns out that \(n\)-convexity with \(n \geq 3\) and convexity are equivalent on \(D_{LIN}\). 2-convex previsions display some drawbacks: in Section 3.2, it is shown that a 2-convex natural extension may be defined and its properties are discussed, but its finiteness is not guaranteed. Moreover, as detailed in Section 3.3, the property of internality may fail (with some limitations, for instance lack of internality cannot be two-sided); agreement with conditional implication (the Goodman-Nguyen relation) is not guaranteed either. In Section 3.4, we show that the special subset of centered 2-convex previsions is not affected by these problems. In Section 4, 2-coherent lower previsions are discussed and characterised on \(D_{LIN}\) (Proposition 9). We compare 2-coherence and \(n\)-coherence in Section 4.1: again, \(n\)-coherence \((n \geq 3)\) and coherence are equivalent on \(D_{LIN}\). \(N\)-coherent previsions \((n \geq 3)\) defined on a generic set of gambles \(S\) have no \(n\)-coherent extension on sufficiently large supersets whenever the equivalence does not hold already on \(S\). We show also that 2-coherence should be preferred to 2-convexity when positive homogeneity and conjugacy are required. The 2-coherent natural extension is introduced and studied in Section 4.2. 2-coherent lower previsions always have it. The extent of the Generalised Bayes Rule for 2-coherent lower previsions is discussed in Section 4.3. Models that can be accommodated into the framework of 2-convexity or 2-coherence, but not of coherence, are presented in Section 5. We focus on how 2-convexity can motivate defining conditional versions of capacities andniveloids, and on the consistency properties of Value-at-Risk, a well-known risk measure which is centered 2-convex, but may even fail to be 2-coherent. In Section 6 we analyse 2-convexity and 2-coherence in a desirability approach. Generalising prior work by Williams [22, 23] for coherence, we focus on the correspondence between these previsions and sets of desirable gambles, and on establishing the ensuing desirability rules. The major differences with the rules for Williams coherence are pointed out in the comments following Propositions 17 and 20. Section 7 concludes the paper. An earlier presentation of the topics in this paper, less extended and without proofs, was delivered at the ISIPTA’15 Symposium [16].
2 Preliminaries

The starting points for our investigation are the known consistency concepts of coherent and convex lower conditional prevision [13, 14, 22, 23]. They both refer to an arbitrary (non-empty) set \( \mathcal{D} \) of conditional gambles, that is of conditional bounded random variables. We denote by \( X|B \) a generic conditional gamble, where \( X \) is a gamble and \( B \) is a non-impossible event \( (B \neq \emptyset) \). It is understood here that \( X: \mathcal{P} \to \mathbb{R} \) is defined on an underlying partition \( \mathcal{P} \) of atomic events \( \omega \), and that \( B \) belongs to the powerset of \( \mathcal{P} \). Therefore, any \( \omega \in \mathcal{P} \) implies either \( B \) or its negation \( \neg B \) (in words, knowing that \( \omega \) is true determines the truth value of \( B \), i.e. \( B \) is known to be either true or false). Given \( B \), the conditional partition \( \mathcal{P}|B \) is formed by the conditional events \( \omega|B \), such that \( \omega \) implies \( B \) (implies that \( B \) is true) and \( X|B : \mathcal{P}|B \to \mathbb{R} \) is such that \( X|B(\omega|B) = X(\omega) \), \( \forall \omega|B \in \mathcal{P}|B \). Because of this equality, several computations regarding \( X|B \) can be performed by means of the restriction of \( X \) on \( B \). In particular, it is useful for the sequel to recall that \( \sup(X|B) = \sup_B X = \sup\{X(\omega) : \omega \in \mathcal{P}, \omega \Rightarrow B\} \), and \( \inf(X|B) = \inf_B X = \inf\{X(\omega) : \omega \in \mathcal{P}, \omega \Rightarrow \neg B\} \).

As a special case, letting \( \Omega \) be the sure event, we have that \( X|\Omega = X \) is an unconditional gamble. Further, \( A|B \) is a conditional event if \( A \) is an event (or its indicator \( I_A \) - we shall generally employ the same notation \( A \) for both).

As customary, without further qualifications, a lower prevision \( \mathcal{P} \) is a map from \( \mathcal{D} \) into the real line, \( \mathcal{P} : \mathcal{D} \to \mathbb{R} \). However, a lower prevision is often interpreted as a supremum buying price [21]. For instance, if a subject assigns \( \mathcal{P}(X|B) \) to \( X|B \), he is willing to buy \( X \), conditional on \( B \) occurring, at any price lower than \( \mathcal{P}(X|B) \). Referring to this behavioural interpretation, the following Definitions 1, 3, 5 require different degrees of consistency for \( \mathcal{P} \), according to whether certain gains depending on \( \mathcal{P} \) avoid losses bounded away from 0. They differ as to the buying and selling constraints they impose.

**Definition 1.** Let \( \mathcal{P} : \mathcal{D} \to \mathbb{R} \) be given.

a) \( \mathcal{P} \) is a coherent conditional lower prevision on \( \mathcal{D} \) iff, for all \( m \in N_0 \), \( \forall X_0|B_0, \ldots, X_m|B_m \in \mathcal{D}, \forall s_0, \ldots, s_m \geq 0 \), defining \( S(s) = \sum s_i B_i(X_i - \mathcal{P}(X_i|B_i)) - s_0 B_0(X_0 - \mathcal{P}(X_0|B_0)) \), it holds, whenever \( S(s) \neq \emptyset \), that \( \sup(S(s)) \geq 0 \).

b) \( \mathcal{P} \) is a convex conditional lower prevision on \( \mathcal{D} \) iff, for all \( m \in N^+ \), \( \forall X_0|B_0, \ldots, X_m|B_m \in \mathcal{D}, \forall s_0, \ldots, s_m \geq 0 \) such that \( \sum s_i = 1 \) (convexity constraint), defining \( G_s = \sum s_i B_i(X_i - \mathcal{P}(X_i|B_i)) - B_0(X_0 - \mathcal{P}(X_0|B_0)) \), \( S(s) = \sum \{B_i : s_i \neq 0, i = 1, \ldots, m\} \), it is sup \( G_s \geq 0 \).

b1) \( \mathcal{P} \) is centered convex or C-convex on \( \mathcal{D} \) iff it is convex and, \( \forall X|B \in \mathcal{D} \), it is \( 0|B \in \mathcal{D} \) and \( \mathcal{P}(0|B) = 0 \).

In the behavioural interpretation recalled above, Definition 1a) considers buying at most \( m \) conditional gambles \( X_1|B_1, \ldots, X_m|B_m \) (also no one, when...
\( m = 0 \) at prices \( P(X_1|B_1), \ldots, P(X_m|B_m) \), respectively, and selling at most one gamble \( X_0|B_0 \) at its supremum buying price \( P(X_0|B_0) \). The gain \( G \) is a linear combination with stakes \( s_0, \ldots, s_m \) of the \( m + 1 \) gains from these transactions. It is conditioned on \( S(s) \), to rule out both trivial transactions \( (G = 0, \text{since } s_0 = \ldots = s_m = 0) \) and the case that \( G = 0 \) because no transaction takes place (when \( B_0, \ldots, B_m \) are all false). Then, coherence requires the non-negativity of the supremum of \( G \), conditional on at least one non-trivial transaction being effective. The interpretation of Definition 1b) is similar: what changes is the convexity constraint on the stakes \( s_0, \ldots, s_m \). This implies that \( G \) is the gain from one selling transaction and at least one buying transaction.

The definition of coherent conditional lower prevision is a structure free version of Williams coherence, discussed in [14]. It is more general than Walley’s coherence [21], in particular it is not necessarily conglomerable and always allows for a natural extension, i.e. there exists an extension on any set of a Williams coherent assessment that is Williams coherent too and least committal. The notion of convex lower prevision is still more general, and was introduced in [13], extending the unconditional convexity studied in [11]. Convex previsions can incorporate various uncertainty models, including convex risk measures, non-normalised possibility measures, and others. However, the special subclass of C-convex lower previsions guarantees better consistency properties. Among these, there always exists a convex natural extension of these measures, whose properties are analogous to those of the natural extension [13, Theorem 9].

Even though coherent and convex lower previsions can be defined on any set of conditional gambles, they are characterised by a few axioms on the special environment \( D_{\text{LIN}} \) defined next.

**Definition 2.** Let \( \mathcal{X} \) be a linear space of gambles and \( \mathcal{B} \subset \mathcal{X} \) a set of (indicators of) events in \( \mathcal{X} \). Suppose that \( \Omega \in \mathcal{B} \) and that \( \mathcal{X} \) is stable by restriction, i.e. \( BX \in \mathcal{X}, \forall B \in \mathcal{B}, \forall X \in \mathcal{X} \). Setting \( B^\emptyset = B - \{\emptyset\} \), define

\[
D_{\text{LIN}} = \{X|B : X \in \mathcal{X}, B \in B^\emptyset\}.
\]  

Note that, since \( \mathcal{B} \subset \mathcal{X} \), the condition \( \Omega \in \mathcal{B} \) implies \( 1 \in \mathcal{X} \) and, therefore, \( \mathcal{X} \) contains all real constants.

The sets \( D_{\text{LIN}} \) may be viewed as generalisations to conditional gambles of linear spaces of unconditional gambles, to which they reduce when \( \mathcal{B} = \{\Omega, \emptyset\} \). Not surprisingly then, characterisations on \( D_{\text{LIN}} \) have an unconditional counterpart on linear spaces.

**Proposition 1.** Let \( P : D_{\text{LIN}} \to \mathbb{R} \) be a conditional lower prevision.

a) \( P \) is coherent on \( D_{\text{LIN}} \) if and only if [23]

(A1) \( P(X|B) - P(Y|B) \leq \sup\{X - Y|B\}, \forall X|B, Y|B \in D_{\text{LIN}}. \)

(A2) \( P(\lambda X|B) = \lambda P(X|B), \forall X|B \in D_{\text{LIN}}, \forall \lambda \geq 0. \)

(A3) \( P(X + Y|B) \geq P(X|B) + P(Y|B), \forall X|B, Y|B \in D_{\text{LIN}}. \)

(A4) \( P(A(X - P(X|A \land B)))|B) = 0, \forall X \in \mathcal{X}, \forall A, B \in B^\emptyset : A \land B \neq \emptyset. \)
b) \( P \) is convex on \( \mathcal{D}_{LIN} \) if and only if (A1), (A4) and the following axiom hold \cite[Theorem 8]{13}:

\[
(A5) \quad P(\lambda X + (1 - \lambda)Y|B) \geq \lambda P(X|B) + (1 - \lambda)P(Y|B), \quad \forall X, Y, B \in \mathcal{D}_{LIN}, \forall \lambda \in [0, 1].
\]

\textbf{Remark 1.} Exploiting some equivalences between axioms or groups of axioms, Proposition 1 as well as the later Propositions 2 and 9 could be restated in a different form. For instance, axiom (A1) is equivalent to the following

\[
(A1') \quad \text{If } X, Y, B \in \mathcal{D}_{LIN}, \mu \in \mathbb{R} \text{ are such that } X|B \geq Y|B + \mu, \text{ then } P(X|B) \geq P(Y|B) + \mu.
\]

Axiom (A1') is also equivalent to monotonicity plus translation invariance:

- If \( X, Y, B \in \mathcal{D}_{LIN} \) and \( X|B \geq Y|B \), then \( P(X|B) \geq P(Y|B) \) (monotonicity).
- If \( X, B \in \mathcal{D}_{LIN} \), \( \mu \in \mathbb{R} \), then \( P(X + \mu|B) = P(X|B) + \mu \) (translation invariance).

Alternatively, (A1) may be replaced in Proposition 1 by \( P(X|B) \geq \inf(X|B) \), \( \forall X, B \in \mathcal{D}_{LIN} \), thus corresponding to the original version in \cite{23}.

Condition (A4) is the Generalised Bayes Rule (GBR), introduced in \cite{22, 23} and studied also in \cite{21} in the special case \( B = \Omega \).

Since our discussion will focus on minimal consistency properties for a conditional lower prevision, we have to mention a generalisation to a conditional framework of the implication (inclusion) relation between events, which is termed Goodman–Nguyen relation \( (\leq_{GN}) \). In fact, suppose \( A \Rightarrow B \) (or \( A \subseteq B \)). Then, asking that \( \mu(A) \leq \mu(B) \) is a really minimal rationality requirement for any \( \mu \) aiming at measuring how likely an event is, given that, whenever event \( A \) proves to be true, \( B \) comes true too. The following extension of the implication to conditional events was proposed in \cite{10}:

\[
A|B \leq_{GN} C|D \quad \text{iff } A \land B \Rightarrow C \land D \quad \text{and } \neg C \land D \Rightarrow \neg A \land B. \quad (2)
\]

The Goodman-Nguyen relation \( \leq_{GN} \) was further extended to conditional gambles in \cite{15}:

\[
X|B \leq_{GN} Y|D \quad \text{iff } I_B X + I_{\neg B \vee D} \sup(X|B) \leq I_D Y + I_{B \vee \neg D} \inf(Y|D)
\]

showing that \( X|B \leq_{GN} Y|D \) implies \( P(X|B) \leq P(Y|D) \) for a C-convex or coherent \( P \) \cite[Proposition 10]{15}.

\textsuperscript{1}Recall that the lower prevision \( P \) is termed convex referring to the convexity constraint \( \sum_{i=1}^{m} s_i = 1 \) in Definition 1 b), not to axiom (A5), which actually tells us that \( P \) is concave, as a real functional.
3 2-convex lower previsions

In Definition 1, a) and b), there is no upper bound to \( m \in \mathbb{N} \). One may think of introducing it as a natural way of weakening coherence and convexity. More precisely, let us call elementary gain on \( X_i|B_i \) any term \( s_iB_i(X_i - P(X_i|B_i)) \), with the proviso that \(-B_0(X_0 - P(X_0|B_0)) \) in Definition 1 b) is also an elementary gain, formally corresponding to \( s_0 = -1 \). Then, we may decide that no more than \( n \) elementary gains are allowed in either \( G \) (Definition 1, a)) or \( G_c \) (Definition 1, b)). When doing so, we speak of \( n \)-coherent or \( n \)-convex lower previsions. This approach extends the notion of \( n \)-coherent (unconditional) prevision in [21, Appendix B].

Intuition suggests that the smaller \( n \) is, the looser the corresponding consistency concept is. In the extreme cases \( n \) may be as small as 1 with coherence, 2 with convexity.

However, 1-coherence is too weak. In fact, \( P \) is 1-coherent on \( \mathcal{D} \) iff, \( \forall X_0|B_0 \in \mathcal{D}, \forall s_0 \in \mathbb{R}, \sup \{ s_0B_0(X_0 - P(X_0|B_0))|B_0 \} \geq 0 \).

It is easy to see that this is equivalent to internality, i.e. to requiring that \( P(X_0|B_0) \in [\inf(X_0|B_0),\sup(X_0|B_0)], \forall X_0|B_0 \in \mathcal{D} \).

**Remark 2.** (1-Avoiding Uniform Loss (1-AUL))

A still weaker concept is that of 1-Avoiding Uniform Loss (1-AUL). Say that \( P \) is 1-AUL on \( \mathcal{D} \) iff \( \forall X_0|B_0 \in \mathcal{D}, \forall s_0 > 0, \)

\[
\sup \{ s_0B_0(X_0 - P(X_0|B_0))|B_0 \} \geq 0. \tag{3}
\]

The condition of 1-AUL is equivalent to \( P(X_0|B_0) \leq \sup(X_0|B_0), \forall X_0|B_0 \in \mathcal{D} \). In particular, this implies \( P(0|B) \leq 0, \forall 0|B \in \mathcal{D} \).

The wording 1-AUL suggests its derivation from a concept of Avoiding Uniform Loss (AUL), which may in fact be obtained from Definition 1 a) by replacing ‘\( m \in \mathbb{N}_0 \)’ with ‘\( m \in \mathbb{N}^+ \)’ (and consequently the gain \( G \) with \( G_{AUL} = \sum_{i=1}^{m} s_iB_i(X_i - P(X_i|B_i)), s_i \geq 0 \).

The notion of 1-AUL has an ancillary role in the theory of 2-convex and 2-coherent lower previsions: it is a rather mild prerequisite to certain properties. In this sense, there is a similarity with the role of the condition of AUL for convex previsions [11, 13] (cf. also Remark 3 in Section 4.2).

Since internality alone does not seem enough as a rationality requirement, we turn our attention in this section to what seems to be the next weakest consistency notion, that is 2-convexity.\(^2\)

**Definition 3.** \( P : \mathcal{D} \rightarrow \mathbb{R} \) is a 2-convex conditional lower prevision on \( \mathcal{D} \) iff, \( \forall X_0|B_0, X_1|B_1 \in \mathcal{D}, \) we have that, defining \( G_{2c} = B_1(X_1 - P(X_1|B_1)) - B_0(X_0 - P(X_0|B_0)) \),

\[
\sup(G_{2c}|B_0 \lor B_1) \geq 0. \tag{4}
\]

\(^2\) 2-convex previsions were termed 1-convex in [2, 15]. Here we prefer the locution ‘2-convex’ by analogy with the rule for fixing \( n \) (as the number of elementary gains) in ‘\( n \)-coherent’ in [21].
3.1 Basic properties of 2-convex lower previsions

We explore now some basic features of 2-convex previsions. Some critical aspects are discussed next, showing in Section 3.4 that they can be solved resorting to the subclass of centered 2-convex previsions.

A remarkable result in our framework is the characterisation of 2-convexity on a structured set $D_{LIN}$.

**Proposition 2.** A conditional lower prevision $P: D_{LIN} \to \mathbb{R}$ is 2-convex on $D_{LIN}$ if and only if (A1) and (A4) hold.

**Proof.** Suppose first that (A1) and (A4) hold. Then, for all $X_0|B_0, X_1|B_1 \in D_{LIN}$, we obtain, using (A1) at the first inequality and (A4) at the second,

$$\sup \{ B_1(X_1 - P(X_1|B_1)) - B_0(X_0 - P(X_0|B_0)) | B_0 \lor B_1 \} \geq P(B_1(X_1 - P(X_1|B_1))|B_0 \lor B_1) - P(B_0(X_0 - P(X_0|B_0))|B_0 \lor B_1) = 0.$$ 

Therefore, $P$ is 2-convex.

Conversely, let $P$ be 2-convex. Then, the proof that (A1) and (A4) hold is part of the proof of Theorem 8 in [13].

To point out an important consequence of Proposition 2, compare it with Proposition 1 b). It follows at once that the difference between 2-convexity and convexity, on $D_{LIN}$, is due to axiom (A5). On the other hand, the proof that a convex prevision on $D_{LIN}$ must satisfy (A5), given in [13, Theorem 8], only involves a gain $G_c$ made up of 3 elementary gains, i.e. it does not fully exploit convexity, but only 3-convexity. This justifies the following conclusion:

**Proposition 3.** On $D_{LIN}$, $n$-convexity with $n \geq 3$ and convexity are equivalent concepts.

Hence, the very difference between convexity and $n$-convexity reduces to that between convexity and 2-convexity, at least on $D_{LIN}$. Yet, if $P$ is defined on a set $D$ other than $D_{LIN}$, we may think of extending it to some $D_{LIN} \supset D$. If $P$ is $n$-convex on $D$, $n \geq 3$, and has an $n$-convex extension to $D_{LIN}$, then $P$ is convex on $D_{LIN}$ and therefore also on $D$. It ensues that if $P$ is $n$-convex ($n \geq 3$) but not convex on $D$, $P$ will have no $n$-convex extension on any sufficiently large superset of $D$ (any $D'$ including some $D_{LIN}$ containing $D$) - see also the later Example 2 in Section 4.1. This is a negative aspect of $n$-convexity, when $n \geq 3$. More generally, the discussion above shows that $n$-convex previsions are not particularly significant as an autonomous concept, when $n \geq 3$.

3.2 The 2-convex natural extension

Turning again to 2-convex previsions, let us define a special extension, the 2-convex natural extension.
Definition 4. Given a lower prevision \( P : \mathcal{D} \to \mathbb{R} \) and an arbitrary conditional gamble \( Z|B \), let

\[
L(Z|B) = \{ \alpha : \sup\{ A(X - P(X|A)) - B(Z - \alpha)|A \lor B \} < 0, \quad \text{for some } X|A \in \mathcal{D} \}.
\]  

(5)

Then the 2-convex natural extension \( E_{2c} \) of \( P \) on \( Z|B \) is

\[
E_{2c}(Z|B) = \sup L(Z|B).
\]  

(6)

In general, \( L(Z|B) \) may be empty, in which case \( E_{2c}(Z|B) = -\infty \), following the usual convention for suprema. When \( L(Z|B) \neq \emptyset \), it is instead possible that \( E_{2c}(Z|B) = +\infty \). The results in the next proposition are helpful in hedging these two occurrences.

Proposition 4. a) \( L(Z|B) \neq \emptyset \), if \( \exists Y|C \in \mathcal{D} \) such that \( C \Rightarrow B \).

b) Let \( P \) be 2-convex and such that \( 0|B \in \mathcal{D} \) and \( P(0|B) = 0 \) \( \forall X|B \in \mathcal{D} \). Given \( 0|C \notin \mathcal{D} \), the extension of \( P \) on \( \mathcal{D} \cup \{0|C\} \) such that \( P(0|C) = 0 \) is 2-convex.

c) When \( L(Z|B) \neq \emptyset \), \( L(Z|B) = \bigcup -\infty, E_{2c}(Z|B) \).

d) If \( L(Z|B) \neq \emptyset \) and \( \sup(X|A) \geq P(X|A) \), \( \forall X|A \in \mathcal{D} \), then \( E_{2c}(Z|B) \leq \sup(Z|B), \forall Z|B \).

e) Let \( P \) be 2-convex and \( 0|B \in \mathcal{D} \), \( \forall X|B \in \mathcal{D} \). Then, \( \forall X|B \in \mathcal{D} \), \( \sup(X|B) \geq P(X|B) \iff P(0|B) \leq 0 \).

Proof of a). Identical to the proof of Proposition 6 in [13].

Proof of b). To check that the extension on \( \mathcal{D} \cup \{0|C\} \) with \( P(0|C) = 0 \) is 2-convex, we only have to check the supreme of two non-trivial gains in Definition 3: the one arising from buying \( X|B \) and selling \( 0|C \), and that corresponding to buying \( 0|C \) and selling \( X|B \).

In the former situation, the gain is

\[
G_{2c} = B(X - P(X|B)) - C(0 - P(0|C)) = B(X - P(X|B)) - B(0 - P(0|B)),
\]

and \( \sup(G_{2c}|B \lor C) \geq \sup(G_{2c}|B) = \sup(B(X - P(X|B)) - B(0 - P(0|B))|B) \geq 0 \), using 2-convexity of \( P \) on \( \mathcal{D} \) at the last inequality.

The latter situation can be treated analogously.

Proof of c). We show first that \( L(Z|B) \subseteq \bigcup -\infty, E_{2c}(Z|B) \). Let \( \alpha \in L(Z|B) \) such that \( s = \sup\{ A(X - P(X|A)) - B(Z - \alpha)|A \lor B \} < 0 \) for some \( X|A \in \mathcal{D} \). By (6), \( \alpha \leq E_{2c}(Z|B) \). If \( \text{ex absurdo } \alpha = E_{2c}(Z|B) \), taking \( \delta > 0 \) such that \( s < s + \delta < 0 \), we get \( \sup\{ A(X - P(X|A)) - B(Z - (\alpha + \delta))|A \lor B \} = \sup\{ A(X - P(X|A)) - B(Z - \alpha) + B\delta|A \lor B \} \leq s + \sup\{ B\delta|A \lor B \} \leq s + \delta < 0 \), a contradiction.

Conversely, let \( \alpha \in [-\infty, E_{2c}(Z|B)] \). Then there exists \( \beta \in L(Z|B) : \alpha < \beta \leq E_{2c}(Z|B) \). Further, \( \sup\{ A(X - P(X|A)) - B(Z - \alpha)|A \lor B \} \leq \sup\{ A(X - P(X|A)) - B(Z - \beta)|A \lor B \} < 0 \), which implies \( \alpha \in L(Z|B) \).
Proof of d). We show that sup(Z|B) \notin L(Z|B). Recalling (5), and since
\[ -B(Z - sup(Z|B)) \geq 0, \sup\{A(X - P(X|A)) - B(Z - sup(Z|B))|A \lor B\} \geq \sup\{A(X - P(X|A))|A \lor B\} \geq \sup\{A(X - P(X|A))|A\} = sup(X|A) - P(X|A) \geq 0. \]
This means that sup(Z|B) \notin L(Z|B). By c), sup(Z|B) \geq E_{2c}(Z|B).

Proof of e). If sup(X|B) \geq P(X|B), \forall X|B \in \mathcal{D}$, then in particular $P(0|B) \leq sup(0|B) = 0$.

As for the reverse implication, let $P(0|B) \leq 0$. Since $P$ is 2-convex, it holds that
\[ 0 \leq \sup\{B(X - P(X|B)) - B(0 - P(0|B))|B\} = \sup\{B(X - P(X|B)) + BP(0|B)|B\} \leq \sup\{X|B - P(X|B)\} = sup(X|B) - P(X|B), \text{ that is } P(X|B) \leq sup(X|B). \]

Parts a) and b) of Proposition 4 suggest a simple way to ensure $E_{2c}(Z|B) \neq -\infty$: just add the gamble 0|B to $\mathcal{D}$, putting $P(0|B) = 0$. To guarantee $E_{2c}(Z|B) \neq +\infty$, it is sufficient that any 0|C in $\mathcal{D}$ (or added to $\mathcal{D}$) is given a non-positive lower prevision, by d) and e). Clearly, the simplest and most obvious choice is to put $P(0|C) = 0, \forall 0|C$. This would make $P$ a centered 2-convex lower prevision; in the remainder of this section we do not however rule out the possibility that $P(0|C) \neq 0$ for some 0|C.

The properties of the 2-convex natural extension are very similar to those of the natural extension:

**Proposition 5.** Let $P : \mathcal{D} \to \mathbb{R}$ be a lower prevision, with $\mathcal{D} \subseteq \mathcal{D}_{LIN}$. If $E_{2c}$ is finite on $\mathcal{D}_{LIN}$, then

a) $E_{2c}(X|B) \geq P(X|B), \forall X|B \in \mathcal{D}$.

b) $E_{2c}$ is 2-convex on $\mathcal{D}_{LIN}$.

c) If $P^*$ is 2-convex on $\mathcal{D}_{LIN}$ and $P^*(X|B) \geq P(X|B), \forall X|B \in \mathcal{D}$, then $P^*(X|B) \geq E_{2c}(X|B), \forall X|B \in \mathcal{D}_{LIN}$.

d) $P$ is 2-convex on $\mathcal{D}$ if and only if $E_{2c} = P$ on $\mathcal{D}$.

e) If $P$ is 2-convex on $\mathcal{D}$, $E_{2c}$ is its smallest 2-convex extension on $\mathcal{D}_{LIN}$.

**Proof.** Assumptions a) \(\vdash e\) of Theorem 9 in [13] (regarding properties of the convex natural extension), with some obvious notation changes and simplifications. In the proof of a) replace $E_c$ in [13] with $E_{2c}$. The proof of b) checks that $E_{2c}$ satisfies axioms (A1) and (A4) (called (D1), (D3) in [13]), according to Proposition 2. This is done using some special gains, simplifying those in [13]: the summation of the terms with stakes $s_1, \ldots, s_m$ is replaced by a single term with stake $s_1 = 1$, in agreement with Definition 3. The proofs of c), d) are analogous, while e) follows from c) and d).

In words, the 2-convex natural extension dominates $P$ (by a)), characterises 2-convexity (by d)) and is the least-committal 2-convex extension of $P$ (by b), c), e)).
3.3 Drawbacks of 2-convexity

Being rather weak a consistency concept, 2-convexity may not satisfy a number of properties which necessarily hold for coherent lower previsions. For instance, the positive homogeneity axiom (A2) of Proposition 1, \( P(\lambda X|B) = \lambda P(X|B) \), with \( \lambda \geq 0 \), may not hold, not even weakening it to

\[
P(\lambda X|B) \geq \lambda P(X|B), \forall \lambda \in [0,1].
\]  

(Unconditional versions of (7) hold for centered convex previsions.)

It can instead be shown that

**Proposition 6.** If, given \( \lambda \in \mathbb{R} \), \( P \) is 2-convex on \( \mathcal{D} \supseteq \{X|B, \lambda X|B\} \), then necessarily

\[
\inf\{(\lambda - 1)X|B\} + P(X|B) \leq P(\lambda X|B) \leq \sup\{(\lambda - 1)X|B\} + P(X|B). \tag{8}
\]

**Proof.** To obtain the first inequality, apply Definition 3 with \( X_1|B_1 = X|B \) and \( X_0|B_0 = \lambda X|B \):

\[
\sup\{B(X - P(X|B)) - B(\lambda X - P(\lambda X|B))|B\} \geq 0 \text{ iff } \\
\sup\{(1 - \lambda)X|B\} - P(X|B) + P(\lambda X|B) \geq 0 \text{ iff } \\
P(\lambda X|B) \geq \inf\{(\lambda - 1)X|B\} + P(X|B).
\]

The proof of the second inequality is analogous (let \( X_1|B_1 = \lambda X|B \), \( X_0|B_0 = X|B \) in Definition 3).

Condition (8) seems to be rather mild, as the next example points out.

**Example 1.** Given \( \mathcal{D} = \{X|B, 2X|B\} \) (\( \lambda = 2 \)), where the image of \( X|B \) is \([-1,1]\) and \( P(X|B) = 0.2 \), equation (8) gives the bounds \( P(2X|B) \in [-0.8,1.2] \). It is easy to check that \( P \) is 2-convex on \( \mathcal{D} \) whatever is the choice for \( P(2X|B) \) in the interval \([-0.8,1.2]\). Depending on the value for \( P(2X|B) \) selected in this interval, it may be \( P(2X|B) \geq 2P(X|B) \).

An annoying feature of 2-convexity is that internality may fail, i.e. \( P(X|B) \) need not belong to the closed interval \( [\inf(X|B), \sup(X|B)] \). Thus, 2-convex previsions may not satisfy a property holding even for 1-coherent previsions.

It has to be noticed that 2-convexity permits no complete freedom in departing from internality. There are two issues to be emphasized with respect to this question. The first tells us that lack of internality cannot be two-sided, because of the following result.

**Proposition 7.** If \( P : \mathcal{D} \rightarrow \mathbb{R} \) is 2-convex on \( \mathcal{D} \) and \( P(Y|D) < \inf(Y|D) \) for some \( Y|D \in \mathcal{D} \), then \( P(X|B) \leq \sup(X|B), \forall X|B \in \mathcal{D} \). Similarly, \( P(Y|D) > \sup(Y|D) \) for some \( Y|D \in \mathcal{D} \) implies \( P(X|B) \geq \inf(X|B), \forall X|B \in \mathcal{D} \).

**Proof.** We equivalently prove that there are no \( X|B, Y|D \in \mathcal{D} \) such that \( P(X|B) > \sup(X|B) \) and \( P(Y|D) < \inf(Y|D) \).
By contradiction, take $\delta, \epsilon > 0$ and suppose

$$P(X|B) = \sup(X|B) + \delta, \quad P(Y|D) = \inf(Y|D) - \epsilon. \quad (9)$$

Then, $G_{2c}|B \lor D = B(X - (\sup(X|B) + \delta)) - D(Y - (\inf(Y|D) - \epsilon))|B \lor D$ is such that $\sup(G_{2c}|B \lor D) < 0$. In fact, $\sup(G_{2c}|B \lor D) = \max\{\sup(G_{2c}|-B \land D), \sup(G_{2c}|B \land \neg D), \sup(G_{2c}|B \lor D)\}$ and we have:

- $\sup(G_{2c}|-B \land D) = \sup(-D(Y - \inf(Y|D) + \epsilon)|B \land D) = \sup(-Y|-B \land D) + \inf(Y|D) - \epsilon = \inf(Y|D) - \inf(Y|-B \land D) - \epsilon \leq -\epsilon < 0$;
- $\sup(G_{2c}|B \land \neg D) = \sup(X|B \land \neg D) - \sup(X|B) - \delta \leq -\delta < 0$;
- $\sup(G_{2c}|B \lor D) = \sup(X|B \lor D) - \sup(X|B) - \delta + \inf(Y|D) - \inf(Y|B \land D) - \epsilon \leq -\delta - \epsilon < 0$.

Therefore, any $P$ satisfying (9) is not 2-convex, according to Definition 3. \qed

The second issue is that 2-convexity imposes a sort of, so to say, two-component internality. To see this, note that

**Lemma 1.** If $P : D \to \mathbb{R}$ is 2-convex on $D$, and $X|B$, $Y|B \in D$, then

$$\inf\{X - Y|B\} \leq P(X|B) - P(Y|B) \leq \sup\{X - Y|B\}. \quad (10)$$

**Proof.** The second inequality in (10) is axiom (A1), a necessary condition for 2-convexity which implies also the first inequality. In fact, $\inf\{X - Y|B\} = -\sup\{Y - X|B\} \leq -(P(Y|B) - P(X|B)) = P(X|B) - P(Y|B)$. \qed

Recall now that $P(X|B)$ is interpreted as a supremum buying price for $X|B$, and that Definition 3 ensures that buying $X|B$ for $P(X|B)$ and selling $Y|B$ at its supremum buying price $P(Y|B)$ would be (marginally) acceptable for 2-convexity. Then, equation (10) tells us that the profit $P(X|B) - P(Y|B)$ from this two-component exchange $(X|B$ vs. $Y|B)$ guarantees no arbitrage. For instance, it cannot exceed the income upper bound $\sup\{X - Y|B\}$.

As a further questionable feature of 2-convexity, the Goodman-Nguyen relation may not induce an agreeing ordering on a 2-convex prevision. This is tantamount to saying that the partial ordering of some 2-convex conditional previsions may conflict with the ordering of the extended implication (inclusion) relation $\leq_{GN}$.

For instance, from (2), if $B \Rightarrow C$ then $0|C \leq_{GN} 0|B$. Agreement with the Goodman-Nguyen relation requires $P(0|C) \leq P(0|B)$ to hold, but it can be proven that if $P(0|B) < 0$ and $B \Rightarrow C$, then 2-convexity asks instead that $P(0|C) \geq P(0|B)$ (the inequality may be strict).
3.4 Centered 2-convex lower previsions

The critical issues of 2-convexity discussed in the preceding section can be solved or softened requiring the additional property

\[ \forall X|B \in \mathcal{D}, 0|B \in \mathcal{D} \text{ and } \mathcal{P}(0|B) = 0, \]

i.e. restricting our attention to centered 2-convex conditional lower previsions. This is shown in the following proposition.

**Proposition 8.** Let \( \mathcal{P} : \mathcal{D} \rightarrow \mathbb{R} \) be a centered 2-convex lower prevision on \( \mathcal{D} \). Then,

a) \( \forall X|B \in \mathcal{D}, \mathcal{P}(X|B) \in [\inf(X|B), \sup(X|B)] \).

b) \( \mathcal{P} \) has a finite 2-convex natural extension \( E_{2c} \) on any superset of \( \mathcal{D} \).

c) \( X|B \leq_{GN} Y|D \) implies \( \mathcal{P}(X|B) \leq \mathcal{P}(Y|D) \).

**Proof.**

*Proof of a).* Put \( Y|B = 0|B \) and \( \mathcal{P}(0|B) = 0 \) in (10).

*Proof of b).* The statement follows from Proposition 4.

*Proof of c).* Proven in [15, Proposition 10] for C-convex previsions. As noted in the Discussion following Proposition 10 in [15], the very same proof applies to centered 2-convex previsions too (cf. also Footnote 2).

**Comment.**

The condition \( \mathcal{P}(0|B) = 0 \) seems to be obvious, and in fact guarantees more satisfactory properties to 2-convexity. In our view, the main reason for considering the alternative \( \mathcal{P}(0|B) \neq 0 \) is to encompass additional uncertainty models. This is patent already in the unconditional framework: convex risk measures, as introduced in [8, 9], correspond to convex, not necessarily centered previsions [11].

Note that centered 2-convexity implies 1-coherence, by Proposition 8 a), while being obviously implied by 2-coherence. Hence, the centering condition \( \mathcal{P}(0|B) = 0 \) may be regarded as a technical instrument to guarantee that the lower prevision \( \mathcal{P} \) ensures more satisfactory properties than a generic 2-convex prevision, without having to assume 2-coherence.

4 2-coherent lower previsions

Our next step is a discussion of which additional properties are achieved by a 2-coherent lower prevision.

**Definition 5.** \( \mathcal{P} : \mathcal{D} \rightarrow \mathbb{R} \) is a 2-coherent lower prevision on \( \mathcal{D} \) iff \( \forall X_0|B_0, X_1|B_1 \in \mathcal{D}, \forall s_1 \geq 0, \forall s_0 \in \mathbb{R}, \) defining \( S(g) = \sqrt{\{B_i : s_i \neq 0, i = 0, 1\}} \), \( G_2 = s_1B_1(X_1 - \mathcal{P}(X_1|B_1)) - s_0B_0(X_0 - \mathcal{P}(X_0|B_0)) \) we have that, whenever \( S(g) \neq \emptyset \),

\[ \sup\{G_2|S(g)\} \geq 0. \]

(11)
2-coherent lower previsions are characterized on $\mathcal{D}_{LIN}$ as follows:

**Proposition 9.** Let $P : \mathcal{D}_{LIN} \to \mathbb{R}$ be a conditional lower prevision. $P$ is 2-coherent on $\mathcal{D}_{LIN}$ if and only if (A1), (A2), (A4) and the following axiom hold:

(A6) $P(\lambda X | B) \leq \lambda P(X | B), \forall \lambda < 0.$

**Proof.** We prove first that if (A1), (A2), (A4) and (A6) hold, then $P$ is 2-coherent on $\mathcal{D}_{LIN}$. Recalling for this Definition 5, take any two $X_0, X_1, B_0, B_1 \in \mathcal{D}_{LIN},$ and any $s_1 \geq 0, s_0 \in \mathbb{R}$. Then, using (A1) at the first inequality, (A2) (when $s_0 \geq 0$) or (A2) and (A6) (when $s_0 < 0$) at the second inequality, we obtain:

$$\sup\{|s_1 B_1(X_1 - P(X_1 | B_1))| - |s_0 B_0(X_0 - P(X_0 | B_0))| |S(\lambda)| \geq P(s_1 B_1(X_1 - P(X_1 | B_1))|S(\lambda)| - P(s_0 B_0(X_0 - P(X_0 | B_0))|S(\lambda)|$$

$$\geq s_1 P(B_1(X_1 - P(X_1 | B_1))|S(\lambda)| - s_0 P(B_0(X_0 - P(X_0 | B_0))|S(\lambda)| = 0,$$

where the equality holds because, when $s_i \neq 0, s_i P(B_i(X_i - P(X_i | B_i))|S(\lambda)) = s_i P(B_i(X_i - P(X_i | B_i) \land S(\lambda))|S(\lambda)) = 0 (i = 1, 2)$ by (A4).

Conversely, if $P$ is 2-coherent, therefore also 2-convex, on $\mathcal{D}_{LIN}$, (A1) and (A4) hold by Proposition 2. Hence, it only remains to prove (A2) and (A6).

We prove first (A6). Apply Definition 5, with $X_1 | B_1 = X | B$, $X_0 | B_0 = \lambda X | B$, $s_1 = 1$, $s_0 = \frac{1}{\lambda} < 0$: sup$(B(X - P(X | B))) - \frac{1}{\lambda} B(\lambda X - P(\lambda X | B))B) = sup(-P(X | B)$$ + \frac{1}{\lambda} P(\lambda X | B)) \geq 0$, which is equivalent to $P(\lambda X | B) \leq \lambda P(X | B)$.

As for (A2), consider the same assumptions of the proof of (A6). Since now $s_0 = \frac{1}{\lambda} > 0$, we obtain the inequality $P(\lambda X | B) \geq \lambda P(X | B)$. Assuming instead $X_1 | B_1 = \lambda X | B$, $X_0 | B_0 = X | B$, $s_1 = \frac{1}{\lambda}, s_0 = 1$, we obtain the reverse inequality $P(\lambda X | B) \leq \lambda P(X | B)$. ☐

**Comment** A comparison of Propositions 2 and 9 is useful for detecting two major differences between (centered) 2-convex and 2-coherent previsions.

One is positive homogeneity (axiom (A2)), a condition which, on any set $\mathcal{D}$, is necessary for 2-coherence, but not for 2-convexity. The need for positive homogeneity depends on the specific model we wish to consider. We might be willing to reject it in some instance, typically because of liquidity risk considerations. Basically, this means that for a large positive $\lambda$ difficulties might be encountered at exchanging $\lambda X | B$ at a price $P(\lambda X | B) = \lambda P(X | B)$, because of lack of market liquidity at some degree.

The second difference is pointed out by axiom (A6). To fix its meaning, recall that given $P(X | B)$, its conjugate upper prevision $\overline{P}(X | B)$ is defined by

$$\overline{P}(X | B) = -P(-X | B).$$

Hence, axiom (A6) ensures by (12) that

$$\overline{P}(X | B) \geq P(X | B), \forall X | B \in \mathcal{D}_{LIN}.$$

Therefore, 2-coherence is preferable to 2-convexity whenever we fix an upper ($\overline{P}$) and a lower ($P$) bound for the uncertainty evaluation of $X | B$, while keeping positive homogeneity.
4.1 2-coherence versus n-coherence

Compare Propositions 9 and 1, a). Recalling that (A6) is a necessary condition for 2-coherence and hence also for coherence, only the superlinearity axiom (A3) distinguishes 2-coherence and coherence on $D_{LIN}$. From this, deductions on the role of n-coherence, $n \geq 3$, can be made which are quite analogue to those on n-convexity in Section 3. This time, it can be shown that any n-coherent lower prevision, $n \geq 3$, must satisfy (A3), and hence that:

**Proposition 10.** On $D_{LIN}$, n-coherence with $n \geq 3$ and coherence are equivalent concepts.

And again, we may in general argue that n-coherence has no special relevance, compared to coherence, when $n \geq 3$. In particular, n-coherent extensions of an n-coherent $P$ exist on sufficiently large sets if and only if $P$ is coherent.

The latter concept is illustrated in the next example, elaborating on Example 2.7.6 in [21].

**Example 2.** Let $IP = \{a, b, c, d\}$ be a partition of the sure event $\Omega$. Define $P$ on the powerset of $IP$ as follows:

- $P(\{\Omega\}) = 1$
- $P(E) = \frac{1}{2}$ if $E$ is made up of 2 or 3 elements of $IP$, one of which is $a$.  
- $P(E) = 0$ otherwise.

It is shown in [21] that $P$ is not coherent, while being 3-coherent, and hence also 3-convex. We show now that $P$ has no 3-convex extension to the linear space $\mathcal{L}(IP)$ of all gambles defined on $IP$.

In fact, suppose a 3-convex extension, also termed $P$, exists, and define $A = a$, $B = a\lor b$, $C = a\lor c$, $D = a\lor d$. Note that, by applying (A1) with $X = \frac{1}{2}A$, $Y = A$ and $B = \Omega$, we get $P(\frac{1}{2}A) \leq P(A) + \sup(-\frac{1}{2}A) = P(A) = 0$. Therefore, also the 3-convex extension of $P$ to $\frac{1}{2}(B + C + D - 1) = \frac{1}{2}A$ should be non-positive. Note also that $P(-1) = -1$ (use (A1) with $X = 0$, $Y = -1$, $B = \Omega$, to get $P(-1) \geq -1$, which is what is needed next; interchanging $X$ and $Y$ in (A1) gives also $P(-1) \leq -1$). By applying axiom (A5) as a necessary condition for 3-convexity, we obtain $P(\frac{1}{2}(B + C + D - 1)) = P(\frac{1}{2}(\frac{1}{2}B + \frac{1}{2}C) + \frac{1}{2}(\frac{1}{2}D - \frac{1}{2})) \geq \frac{1}{2}P(\frac{1}{2}B + \frac{1}{2}C) + \frac{1}{2}P(\frac{1}{2}D - \frac{1}{2}) \geq \frac{1}{2}P(B) + \frac{1}{2}P(C) + \frac{1}{2}P(D) + \frac{1}{2}P(-1) \geq 3 \cdot \frac{1}{2} \cdot \frac{1}{2} - \frac{1}{2} = \frac{1}{8} > 0$, a contradiction.

From what we have just proven, we may conclude that:

a) the given $P$ on the powerset of $IP$ has no 3-convex extension to $\mathcal{L}(IP)$;

b) $P$ (viewed now as 3-coherent on the powerset of $IP$) has no 3-coherent extension on $\mathcal{L}(IP)$ either: if it had one, this extension would be 3-convex too, contradicting a).

It is interesting to realise that, on $D_{LIN}$, convexity is what is missing to 2-coherence (and viceversa) to achieve coherence:
Proposition 11. $\mathcal{P}: \mathcal{D}_{LIN} \rightarrow \mathbb{R}$ is coherent iff it is both 2-coherent and convex.

Proof. Clearly, coherence implies both 2-coherence and convexity.

Conversely, let $\mathcal{P}$ be both 2-coherent and convex on $\mathcal{D}_{LIN}$. By Propositions 1 and 9, it only remains to check that (A3) holds to ensure coherence of $\mathcal{P}$. In fact, by (A5) and (A2), we have

$$
P(X + Y|B) = P(\frac{1}{2}(2X) + \frac{1}{2}(2Y)|B) \geq \frac{1}{2}P(2X|B) + \frac{1}{2}P(2Y|B) = P(X|B) + P(Y|B).
$$

(13)

4.2 The 2-coherent natural extension

2-coherent lower probabilities, being also centered 2-convex and 1-AUL, always have a 2-convex natural extension.

Appreciably, they further ensure the existence of a 2-coherent natural extension. This tells us that the additional properties of 2-coherence are stable, in the sense that they can be preserved by extension to any set of conditional gambles.

The role of the 2-coherent natural extension is analogous for 2-coherence to that of the 2-convex natural extension for 2-convexity, and most derivations are quite similar. We shall demonstrate in detail only the most differing ones.

Definition 6. Given a lower prevision $\mathcal{P}: \mathcal{D} \rightarrow \mathbb{R}$ and an arbitrary $Z|B$, let

$$
L_2(Z|B) = \{ \alpha : \sup\{s_1 A(X - \mathcal{P}(X|A)) - B(Z - \alpha)|S(s)\} < 0
$$

for some $X|A \in \mathcal{D}, s_1 \geq 0\}.

where

$$
S(s) = \begin{cases} 
A \lor B & \text{if } s_1 > 0 \\
B & \text{if } s_1 = 0.
\end{cases}
$$

Then, the 2-coherent natural extension $E_2$ of $\mathcal{P}$ on $Z|B$ is

$$
E_2(Z|B) = \sup L_2(Z|B).
$$

Proposition 12 (Existence of the 2-coherent natural extension). Given a lower prevision $\mathcal{P}: \mathcal{D} \rightarrow \mathbb{R}$,

a) $L_2(Z|B)$ is non-empty, $\forall Z|B$.

b) $L_2(Z|B) = | - \infty, E_2(Z|B)|$.

c) If $\mathcal{P}$ is 1-AUL, $E_2(Z|B) \leq \sup(Z|B), \forall Z|B$.

d) If $\mathcal{P}$ is not a 1-AUL prevision, $\exists Z|B \in \mathcal{D}$ such that $E_2(Z|B) = +\infty$. 

16
Proof. Proof of a). Take $\alpha < \inf(Z|B)$ and let $X|A$ be any conditional gamble in $D$. Then, putting $s_1 = 0$ in Definition 6, $\sup\{s_1A(X - P(X|A)) - B(Z - \alpha)|S(s)\} = \sup\{-B(Z - \alpha)|B\} = \alpha + \sup\{-Z|B\} = \alpha - \inf\{Z|B\} < 0$, i.e. $\alpha \in L_2(Z|B)$.

Proofs of b) and c). Same as the proofs of, respectively, Proposition 4c) and Proposition 4d), replacing $A(X - P(X|A))$ with $s_1A(X - P(X|A))$.

Proof of d). Let $P$ be not 1-AUL. Then, $\exists Z|B \in D, s > 0$ such that $\sup\{sB(Z - P(Z|B)|B\} < 0$. Clearly, this implies $P(Z|B) - \sup\{Z|B\} > 0$.

Then, $\forall s_1 > 1, \forall \alpha$ such that $\alpha < \sup\{Z|B\} + s_1(P(Z|B) - \sup\{Z|B\})$, we get $\sup\{s_1B(Z - P(Z|B)) - B(Z - \alpha)|B\} = \alpha - \sup\{Z|B\} - s_1(P(Z|B) - \sup\{Z|B\}) < 0$, i.e. $\alpha \in L_2(Z|B)$. Since $\alpha$ can be chosen arbitrarily large in $L_2(Z|B)$ by increasing $s_1, E_2(Z|B) = +\infty$.

□

Remark 3. Proposition 12 ensures that the condition of 1-AUL is necessary and sufficient for the finiteness of the 2-coherent natural extension $E_2$.

Proposition 13 (Properties of the 2-coherent natural extension). Let $D \subseteq D_{LIN}$ and $P : D \to \mathbb{R}$. If $E_2$ is finite on $D_{LIN}$, then

a) $E_2(X|B) \geq P(X|B), \forall X|B \in D$.

b) $E_2$ is 2-coherent on $D_{LIN}$.

c) If $P'$ is 2-coherent on $D_{LIN}$ and $P'(X|B) \geq P(X|B), \forall X|B \in D$, then $P' \geq E_2$ on $D_{LIN}$.

d) $P$ is 2-coherent on $D$ if and only if $E_2 = P$ on $D$.

e) If $P$ is 2-coherent on $D$, $E_2$ is its smallest 2-coherent extension on $D_{LIN}$.

Proof. Proof of a). See the proof of Theorem 9, a) in [13].

Proof of b). By Proposition 9, we prove that $E_2$ satisfies axioms (A1), (A2), (A4), (A6).

(A1) and (A4): see the proof that, respectively, axioms (D1) and (D3) hold in [13], Theorem 9, b) (with the obvious modifications $G_1 = s_1B_1(X_1 - P(X_1|B_1)), s_1 \geq 0$).

(A2): The proof corresponds to that for axiom (A2) in [14], Theorem 3, p. 625 (conditioning now $G_1$ on $S(s)$).

(A6): To prove (A6), follow the next two steps:

1) It should be proven that

$$E_2(\lambda X|B) \leq \lambda E_2(X|B), \forall \lambda < 0.$$  

Since we already know that (A2) holds, we can use it to write $E_2(\lambda X|B) = E_2(-\lambda(-X)|B) = -\lambda E_2(-X|B) \leq \lambda E_2(X|B)$ if and only if $-E_2(-X|B) \geq E_2(X|B)$.

Therefore, we can equivalently prove instead that

$$E_2(X|B) + E_2(-X|B) \leq 0.$$  (14)
2) We prove (14). Take arbitrarily $\alpha^+ \in L_2(X|B)$, $\alpha^- \in L_2(-X|B)$. Letting $g_i = s_iB_i(X_i - P(X_i|B_i))$, $s_i \geq 0$ for $i = 1, 2$, we have by the definition of $L_2(X|B)$, $L_2(-X|B)$:

$$\sup(g_1 - B(X - \alpha^+)|B_1 \lor B) = \sup(Z_1|B_1 \lor B) < 0,$$

$$\sup(g_2 - B(-X - \alpha^-)|B_2 \lor B) = \sup(Z_2|B_2 \lor B) < 0.$$ 

Defining $H = B \lor B_1 \lor B_2$, it holds also that

$$\sup(Z_1 + Z_2|H) = \sup(g_1 + g_2 + B(\alpha^+ + \alpha^-)|H) < 0. \quad (15)$$

In fact, decompose $H$ into the sum of 4 disjoint events as follows:

$$H = B \lor (B_1 \land B_2 \land \neg B) \lor (\neg B_1 \land B_2 \land \neg B) \lor (B_1 \land \neg B_2 \land \neg B). \quad (16)$$

Now condition $Z_1$ and $Z_2$ on any of the 4 events in (16) that are not impossible. Let $H_j$ be the generic such event and $J = \{ j \in \{1, 2, 3, 4\} : H_j \neq \emptyset \}$. Then note that $\sup\{Z_i|H_j\} \leq 0, (i = 1, 2)$. In fact, considering $Z_1$, if $H_j$ is any of $B$, $B_1 \land B_2 \land \neg B$ or $B_1 \land \neg B_2 \land \neg B$, then $\sup\{Z_i|H_j\} \leq \sup\{Z_1|B_1 \lor B\} < 0$, whilst $\sup\{Z_1|\neg B_1 \land B_2 \land \neg B\} = \sup\{0|\neg B_1 \land B_2 \land \neg B\} = 0$. Similarly, $\sup\{Z_2|H_j\} \leq 0$, with equality iff $H_j = B_1 \land \neg B_2 \land \neg B$. It ensues also that the two suprema $\sup\{Z_i|H_j\}$, $\sup\{Z_2|H_j\}$ cannot be simultaneously null, for $j \in J$.

Hence, $\sup(Z_1|H) + \sup(Z_2|H) = \max\{\sup\{Z_1|H_j\} + \sup\{Z_2|H_j\}, j \in J\} < 0$. The inequality (15) follows, since $\sup(Z_1 + Z_2|H) \leq \sup(Z_1|H) + \sup(Z_2|H)$. Further,

$$\sup(g_1 + g_2|H) \geq \sup(g_1 + g_2|B_1 \lor B_2) \geq 0.$$ 

using 2-coherence of $P$ on $D$ at the last inequality.

Therefore, in order for inequality (15) to hold, necessarily $\alpha^+ + \alpha^- < 0$, i.e. $\alpha^- < -\alpha^-, \forall \alpha^+ \in L_2(X|B), \alpha^- \in L_2(-X|B)$. Equivalently,

$$\sup\{\alpha^+ \in L_2(X|B)\} = E_2(X|B) \leq \inf\{-\alpha^- : \alpha^- \in L_2(-X|B)\} = -\sup\{\alpha^- \in L_2(-X|B)\} = -E_2(X|B),$$

which gives (14).

Proofs of c) and d). Analogous to the proof of Theorem 9, c) and d) in [13].

Proof of e). Implied by c) and d).

We may thus conclude that centered 2-convexity and 2-coherence appear to be the most significant weakenings of (centered) convexity and coherence.

4.3 About the Generalised Bayes Rule

By Propositions 2 and 9, the Generalised Bayes Rule (GBR) is a necessary consistency condition for both 2-convex and 2-coherent lower previsions. This
guarantees that this key updating rule holds even with weaker consistency concepts than coherence or convexity. However, it would be erroneous to believe that nothing about the GBR changes with such looser consistency requirements. To see this, put \( B = \Omega \) in (A4), getting
\[
P(A(X - P(X|A))) = 0,
\]
which informs us that \( P(X|A) \) is a solution of the equation
\[
P(A(X - r)) = 0. \tag{17}
\]
From Proposition 9 in [13], we know that if \( P \) is convex on \( \mathcal{D} \supset \{ A, X|A, A(X - P(X|A)) \} \) and \( P(A) > 0 \), then \( P(X|A) \) is the unique solution of (17). A uniqueness result for coherent lower previsions is given in [21], Sec. 6.4.1.

With 2-coherent or 2-convex lower previsions, \( P(X|A) \) may no longer be the unique solution of (17). The next result illustrates this for 2-coherence.

**Proposition 14.** Let \( P : \mathcal{D} = \{ A, X|A, A(X - r), A(X - q) \} \rightarrow \mathbb{R} \) be a lower prevision, such that \( r \neq q \), \( P(A(X - r)) = P(A(X - q)) = 0 \), \( A \neq \Omega \), and \( 1 \geq P(A) > 0 \). Then \( P \) is 2-coherent on \( \mathcal{D} \) if and only if
\[
P(X|A), r, q \in [\inf(X|A), \sup(X|A)]. \tag{18}
\]

**Proof.** Suppose first that (18) holds. To prove that \( P \) is 2-coherent on \( \mathcal{D} \), we may check by Definition 5 that any admissible gain \( G_2 \) satisfies (11). For this, we consider the gains from betting on all couples of elements of \( \mathcal{D} \) (and their special cases \( s_0 = 0 \) or \( s_1 = 0 \), where the effective bet is on a single element). These gains may be partitioned into two groups:

1. Gains from bets on the couples \((X|A, A(X - r)), (X|A, A(X - q)), (A(X - r), A(X - q))\).

The proofs that all such gains satisfy (11) are very similar for all couples. To exemplify, take the couple \((X|A, A(X - r))\). Any admissible gain is either of
\[
G_2 = s_1 A(X - P(X|A)) - s_0 A(X - r); \quad G_2' = s_1 A(X - r) - s_0 A(X - P(X|A)).
\]

Let us first look at \( G_2 \). If \( s_0 \neq 0 \), \( G_2|S(g) = G_2|\Omega \), and \( \sup G_2 \geq G_2(\neg A) = 0 \). If \( s_0 = 0 \) (and \( s_1 > 0 \)), \( \sup(G_2|S(g)) = \sup(G_2|A) = s_1 \sup(X|A - P(X|A)) \geq 0 \) by (18).

Consider now \( G_2' \). If \( s_1 \neq 0 \), \( G_2'|S(g) = G_2'|\Omega \), and \( \sup G_2' \geq G_2'(\neg A) = 0 \). Let then \( s_1 = 0 \), hence \( S(g) = A \).

If \( s_0 > 0 \), \( \sup(G_2'|S(g)) = \sup(-s_0 A(X - P(X|A))|A) = s_0 \sup(P(X|A) - X|A) = s_0 (P(X|A) - \inf(X|A)) \geq 0 \) by (18).

If \( s_0 < 0 \), \( \sup(G_2'|S(g)) = \sup(-s_0 A(X - P(X|A))|A) = -s_0 (\sup(X|A) - P(X|A)) \geq 0 \), again by (18).
2. Gains from betting on one of the remaining three couples. All such couples include $A$, and the proof is identical for each of them. Take for instance the couple $(A, X|A)$. The related gains are

$$G_2 = s_1(A - P(A)) - s_0(A - P(X|A));$$
$$G'_2 = s_1A(X - P(X|A)) - s_0(A - P(A)).$$

Consider $G_2$. If $s_1 = 0$ (and $s_0 \neq 0$), $G_2$ coincides with $G'_2$ in 1., case $s_1 = 0$. Hence the same derivation and conclusions apply.

Let now $s_1 > 0$, hence $S(\bar{g}) = \Omega$. Then $\sup(G_2|\Omega) \geq \sup(G_2|A) \geq 0$. The last inequality holds because

$$\sup(G_2|A) = s_1(1 - P(A)) + s_0P(X|A) + \sup(-s_0X|A) \quad (19)$$

and from (18), (19) we obtain:

- if $s_0 \geq 0$, $\sup(G_2|A) \geq s_1(1 - P(A)) + s_0\inf(X|A) - s_0\inf(X|A) \geq 0$;
- if $s_0 < 0$, $\sup(G_2|A) \geq s_1(1 - P(A)) + s_0\sup(X|A) - s_0\sup(X|A) \geq 0$.

Referring to $G'_2$, if $s_0 = 0$ then $S(\bar{g}) = A$, $\sup(G'_2|A) = s_1(\sup(X|A) - P(X|A)) \geq 0$ by (18).

When $s_0 \neq 0$, $S(\bar{g}) = \Omega$ and sup $G'_2 = \max\{\sup(G'_2|A), \sup(G'_2|\neg A)\} \geq 0$. In fact, if $s_0 > 0$ then $\sup(G'_2|\neg A) = s_0P(A) > 0$. If $s_0 < 0$, $\sup(G'_2|A) = -s_1P(X|A) - s_0(1 - P(A)) + s_1\sup(X|A) \geq -s_1\sup(X|A) - s_0(1 - P(A)) + s_1\sup(X|A) \geq 0$.

Conversely, let now $P$ be 2-coherent. Since $P$ is also 1-coherent, $P(X|A)$ satisfies condition (18). To see that also $r$ does so (the proof for $q$ is identical), note that the gain

$$G_2 = s_1(A - P(A)) - A(X - r), s_1 > 0$$

is such that $\sup(G_2) \geq 0$, by (11). Since $\sup(G_2|\neg A) = -s_1P(A) < 0$, necessarily $\sup(G_2|A) = s_1(1 - P(A)) + r - \inf(X|A) \geq 0$, that is

$$\inf(X|A) \leq r + s_1(1 - P(A)), \forall s_1 > 0.$$

From the above inequality, $r \geq \inf(X|A)$.

To prove that $r \leq \sup(X|A)$, consider the gain

$$G'_2 = A(X - r) - s_0(A - P(A)), s_0 < 0,$$

and note that $G'_2|\neg A = s_0P(A) < 0$. This implies, for any $s_0 < 0$, $\sup(G'_2|A) = \sup(X|A) - r - s_0(1 - P(A)) \geq 0$. Hence, $r \leq \sup(X|A)$.

\textit{Comment.} Proposition 14 establishes that equation (17) has more solutions, when $P$ is 2-coherent. Actually, there are infinitely many, provided they comply with the internality condition (18). Lack of uniqueness means also that we are
not obliged to choose one of these solutions: any two of them can 2-coherently coexist in the set $\mathcal{D}$ of Proposition 14. Even as many solutions as we wish may be found in a single 2-coherent assessment. Just think that this does not essentially alter the proof of Proposition 14, since 2-coherence restricts checking it on gains referring to (at most) couples of gambles.

Since a 2-coherent prevision is also 2-convex, it is clear that the GBR will generally not be the unique solution of equation (17) even when $\mathcal{P}$ is 2-convex. We omit detailing this case.

5 Weakly consistent uncertainty models

As mentioned in the Introduction, a motivation for studying the loose forms of consistency introduced in this paper is their capability of encompassing or extending uncertainty models already investigated in the literature. Even though these models may depart also considerably from coherence and convexity, they can nevertheless be accommodated into a unifying betting scheme, ranging from 2-convex to coherent lower previsions.

5.1 Capacities and niveloids

Focusing on 2-convexity, we first recall a few definitions and some results concerning unconditional 2-convex lower previsions.

Definition 7. Given a finite partition $\mathcal{P}$, and denoting with $2^{\mathcal{P}}$ its powerset, a mapping $c : 2^{\mathcal{P}} \to [0, 1]$ is a (normalised) capacity whenever $c(\emptyset) = 0$, $c(\Omega) = 1$ (normalisation) and $\forall A_1, A_2 \in 2^{\mathcal{P}}$ such that $A_1 \Rightarrow A_2$, $c(A_1) \leq c(A_2)$ (1-monotonicity).

Definition 8. Given a linear space $\mathcal{L}$ of random variables, a niveloid $[3, 7]$ is a functional $N : \mathcal{L} \to \mathbb{R} = \mathbb{R} \cup \{-\infty, +\infty\}$ which is translation invariant and monotone, i.e. such that

$$N(X + \mu) = N(X) + \mu, \forall X \in \mathcal{L}, \forall \mu \in \mathbb{R};$$

$$X \geq Y \text{ implies } N(X) \geq N(Y), \forall X, Y \in \mathcal{L}. \quad (20)$$

As well-known, capacities are uncertainty measures with really minimal quantitative requirements. Niveloids can be viewed as a generalisation of capacities to linear spaces of random variables which preserves their minimality properties. Strictly speaking, this is true for centered niveloids, i.e. such that $N(0) = 0$. In fact, the centering condition $N(0) = 0$ does not ensue from the definition of niveloid. Note also that niveloids apply to random variables which may be unbounded too.

It has been proven in [2, Section 4.1]$^3$ that:

Proposition 15. a) Let $\mathcal{P}$ be defined on $2^{\mathcal{P}}$. Then $\mathcal{P}$ is a centered 2-convex lower prevision if and only if it is a capacity.

$^3$See Footnote 2.
b) Let $P$ be defined on a linear space $L$ of gambles. Then $P$ is a 2-convex lower prevision if and only if it is a (finite-valued) niveloid.

Hence, an unconditional 2-convex lower prevision is equivalent to a capacity or a niveloid, on structured sets ($2^P$ or $L$ respectively). On non-structured sets, it extends these concepts.

2-convex conditional lower previsions are natural candidates to define conditional capacities and niveloids on arbitrary sets of, respectively, conditional events or gambles. To the best of our knowledge, such conditional versions have not been considered yet in this general conditional environment, but rather in more specific cases. For instance, [4] focuses on updating rules for ‘convex’ capacities, which means for 2-monotone lower probabilities, while considering a single conditioning event.

Thus 2-convex previsions may provide an appropriate framework for such extensions, guaranteeing some minimal properties like the existence of a 2-convex natural extension (when being centered). Take for instance centered 2-convex conditional lower probabilities. They satisfy the properties one would require to a conditional capacity: $P(0|B) = 0$, $P(\Omega|B) = 1$ (this follows from Proposition 8, a)), and $A|B \leq_{GN} C|D$ implies $P(A|B) \leq P(C|D)$ (Proposition 8, c)). Similarly, centered 2-convex lower previsions ensure generalisations of properties (20) (see especially Proposition 2 and Remark 1 for the first property, Proposition 8, c) for the second).

5.2 Value-at-Risk (VaR)

Several examples of weakly consistent models may be found among the many risk measures that have been proposed in the financial literature. We shall discuss here Value-at-Risk (VaR), probably the most widespread such measure.

A risk measure $\rho$ is a map $\rho : D \to \mathbb{R}$ assigning a number $\rho(X)$ to each gamble $X \in D$, aiming at measuring how ‘risky’ $X$ is. Risk measures are strongly connected to imprecise previsions: any risk measure $\rho(X)$ on $X$ corresponds to the opposite $-P(X)$ of a lower prevision for $X$ [12].

Because of this correspondence, we may transpose concepts developed for imprecise probability theory to risk measurement (and vice versa). Hence it is possible to check whether a certain risk measure is coherent, convex, or at least 2-coherent or 2-convex, according to whether the corresponding $P = -\rho$ is so.

As for VaR, it is essentially a quantile-based measure:

\begin{definition} ([1]). Given a gamble $X$, a probability $P$ on $\{(X \leq x) : x \in \mathbb{R}\}$, and a real $\alpha \in [0, 1]$, the Value-at-Risk of $X$ at level $\alpha$ is: \[ VaR_\alpha(X) = -\inf\{x \in \mathbb{R} : P(X \leq x) > \alpha\}. \] \end{definition}

\begin{equation} \tag{21} \end{equation}

\footnote{In alternative definitions of VaR, cf. [6, Sec. 2.3.1], the minus in (21) is omitted (this corresponds to reasoning in terms of losses) and/or the strict inequality in the inf is weak. Their consistency properties are the same.}
It is known that VaR is not coherent, although it may be so under some additional, rather strong assumptions [12]. Which are then its guaranteed consistency properties? This amounts to investigating the consistency of a lower prevision \( P^\alpha_V(X) = -\text{VaR}_\alpha(X) \), by the correspondence mentioned above. The next proposition ensures that VaR is centered 2-convex, while Example 3 shows that it may even fail to be 2-coherent.

**Proposition 16.** Let \( L \) be a linear space of gambles, \( \alpha \in [0,1[ \) and \( P \) a probability on \( \bigcup_{X \in L} \{(X \leq x) : x \in \mathbb{R} \} \). Define \( P^\alpha_V \) as

\[
P^\alpha_V(X) = -\text{VaR}_\alpha(X) = \inf\{x \in \mathbb{R} : P(X \leq x) > \alpha\}, \forall X \in L.
\]

Then \( P^\alpha_V \) is a centered 2-convex lower prevision.

**Proof.** By Proposition 15 b), \( P^\alpha_V \) is 2-convex iff it is a niveloid, that is iff it is translation invariant and monotone. Proving translation invariance and monotonicity is essentially the same as proving that \( \text{VaR}_\alpha \) has these properties, which is well known (cf. [6, Sec. 2.3.2]).

As for centering, recalling (21) we have \( P^\alpha_V(0) = -\text{VaR}_\alpha(0) = 0 \).

**Example 3** (VaR may be 2-incoherent). Let \( X \) be a 2-valued gamble such that \( P(X = -1) = P(X = 2) = 0.5 \). Given \( \alpha = 0.6 \), it is easy to check using (21) that \( P^\alpha_V(X) = -\text{VaR}_{0.6}(X) = \inf\{x : P(X \leq x) > 0.6\} = 2 \), while \( P^\alpha_V(-X) = 1 \). Hence \( P^\alpha_V(-X) > -P^\alpha_V(X) \), meaning that axiom (A6) with \( \lambda = -1 \) does not hold for \( P^\alpha_V \). Since (A6) is a necessary condition for 2-coherence, \( P^\alpha_V \) is not 2-coherent.

**Remark 4** (2-coherent models). The models we have seen so far are 2-convex. 2-coherence arises naturally with interval evaluations made up of a lower \( P \) and an upper \( \overline{P} \) uncertainty measure, both 2-convex, like a capacity and its conjugate. In fact, it is then natural to require that \( P \leq \overline{P} \), which is a follow-up of equation (12), implied by 2-coherence. As another instance, we mention \( p \)-boxes. While univariate \( p \)-boxes satisfy stronger consistency properties (they correspond to a couple \((P, \overline{P})\), where both \( P, \overline{P} \) are precise probabilities), bivariate \( p \)-boxes may be related with 2-coherence (cf. [17, Sec. 3.1]).

### 6 Weak consistency in a desirability approach

In this section we examine 2-convexity and 2-coherence from the viewpoint of desirability. This is an alternative approach to rationality concepts for uncertainty measures going back to [22] in the case of conditional imprecise previsions. It has been recently applied to a variety of other situations, see e.g. the discussion in [18] and the results in [19].

Roughly speaking, a set \( \mathcal{A} \) of gambles is considered.\(^5\) It is such that its gambles are regarded as desirable or acceptable. We may in general be willing

\( ^5\)As will appear later, \( \mathcal{A} \) is included into some fixed linear space of gambles.
to establish some rationality criteria, requiring that certain gambles do, or do not, belong to \( \mathcal{A} \). The basic problem we shall consider here is: which is the correspondence between the rationality criteria we adopt and the consistency concepts of centered 2-convexity or alternatively 2-coherence? More specifically, the following two questions arise:

Q1) Which rationality criteria should be required to the elements of a set \( \mathcal{A} \), so that a conditional lower prevision \( \mathcal{P} \) may be obtained from \( \mathcal{A} \) that is 2-coherent (alternatively, 2-convex)?

Q2) Conversely, given a 2-coherent (alternatively, 2-convex) \( \mathcal{P} \), does it determine a set \( \mathcal{A}' \) with certain rationality properties?

In the case that \( \mathcal{P} \) is coherent, the answer to Q1) and Q2) was given by Williams in [22]. Our approach to solving Q1) and Q2) was largely influenced by his work.

Preliminarily, some notation must be introduced.

**Definition 10.** Let \( \mathcal{X} \) be a linear space of gambles, \( \mathcal{B} \subset \mathcal{X} \) a set of (indicators of) events, \( \mathcal{B}^\varnothing = \mathcal{B} \setminus \{\varnothing\} \). We suppose \( \Omega \in \mathcal{B} \) and \( \forall X \in \mathcal{X}, \forall B \in \mathcal{B} \).

Define then

\[
\mathcal{X}^\succeq = \{X \in \mathcal{X} : \inf X \geq 0\}, \\
\mathcal{X}^\preceq = \{X \in \mathcal{X} : \sup X \leq 0\},
\]

and, \( \forall B \in \mathcal{B} \),

\[
\mathcal{R}(B) = \{X \in \mathcal{X} : BX = X\}, \\
\mathcal{R}(B)^\succeq = \{X \in \mathcal{R}(B) : \inf\{X|B\} > 0\}, \\
\mathcal{R}(B)^\preceq = \{X \in \mathcal{R}(B) : \sup\{X|B\} < 0\}.
\]

If \( \mathcal{S} \) and \( \mathcal{T} \) are subsets of \( \mathcal{X} \), their Minkowski sum is

\[
\mathcal{S} + \mathcal{T} = \{X + Y : X \in \mathcal{S}, Y \in \mathcal{T}\}.
\]

We shall use similar compact notation later. For instance, \( \lambda \mathcal{S} + \mu \mathcal{T} \subseteq \mathcal{U} \), \( \forall \lambda, \mu \geq 0 \), means: \( \forall X \in \mathcal{S}, \forall Y \in \mathcal{T}, \forall \lambda, \mu \geq 0, \lambda X + \mu Y \in \mathcal{U} \).

**Lemma 2.** Properties of the sets \( \mathcal{R}(B) \):

a) \( \forall B, \mathcal{R}(B) \) is a linear space.

b) If \( X \in \mathcal{R}(B) \) and \( B \Rightarrow A \), then \( X \in \mathcal{R}(A) \).

**Proof.** a) is trivial. As for b), we have \( AX = I_A X = (I_B + I_{A \wedge \neg B}) X = I_B X = X \).

\[\square\]

\[\text{Note that if } X \in \mathcal{X} \text{ and } B \in \mathcal{B}^\varnothing, X|B \in \mathcal{D}_{L1N} \text{ in the notation of the preceding sections.}\]
6.1 Desirability axioms for 2-coherent previsions

The following proposition answers question Q1) completely for 2-coherence:

**Proposition 17.** Let \( A \subseteq \mathcal{X} \) be such that

a) \( \lambda A + \mathcal{R}(B) \supseteq A, \forall \lambda \geq 0, \forall B \in \mathcal{B}; \)

b) \( \mathcal{R}(B) \cap A = \emptyset, \forall B \in \mathcal{B}. \)

c) \( (\mathcal{R}(B_1) \cap A) + (\mathcal{R}(B_2) \cap A) \subseteq \mathcal{R}(B_1 \lor B_2) \setminus \mathcal{R}(B_1 \lor B_2)^\ominus, \forall B_1, B_2 \in \mathcal{B}. \)

Define, \( \forall X | B \in \mathcal{D}_{LIN}, \)

\[
P(X|B) = \sup \{ x \in \mathbb{R} : B(X - x) \in A \}. \tag{22}
\]

Then, \( P \) is 2-coherent on \( \mathcal{D}_{LIN}. \)

**Proof.** The core idea of the proof is to show that, for any given \( X_0 | B_0, X_1 | B_1 \in \mathcal{D}_{LIN}, \forall s_1 \geq 0, \forall s_0 \in \mathbb{R}, P \) defined by (22) is such that \( G_2 \) satisfies condition (11) in Definition 5, and therefore \( P \) is 2-coherent.

For this, define first the following gambles:

\[
K = \sup \left( \sum_{s_i \neq 0} B_i | S(s) \right), \tag{23}
\]

\[
S_i = \begin{cases} 
|s_i| B_i(X_i - P(X_i|B_i)) + \frac{s_i \epsilon}{|s_i|K} B_i & \text{if } s_i \neq 0 \\
0 & \text{if } s_i = 0 
\end{cases} \tag{24}
\]

In equation (24), \( i = 0, 1 \) and \( \epsilon > 0. \)

Note that \( K > 0, \) as it can take values in \( \{1, 2\}. \) We analyse now the relationships between \( S_0, S_1 \) in (24) and \( A. \) The following facts will be used later in the proof, when expressing \( G_2 \) in terms of \( S_0, S_1 \) and \( K. \)

i) If \( s_i > 0, S_i = s_i B_i(X_i - P(X_i|B_i)) + \frac{s_i \epsilon}{s_i K} B_i \in A, i = 0, 1. \)

In fact, by the definition of \( P \) in (22), \( \exists t \in [0, \frac{s_i}{s_i K}] \) such that \( B_i(X_i - (P(X_i|B_i) - t)) \in A. \)

Writing then

\[
S_i = s_i B_i(X_i - (P(X_i|B_i) - \frac{s_i \epsilon}{s_i K})) = s_i B_i(X_i - (P(X_i|B_i) - t)) + s_i B_i(\frac{s_i \epsilon}{s_i K} - t),
\]

we note that the second term in the summation, \( s_i B_i(\frac{s_i \epsilon}{s_i K} - t), \) belongs to \( \mathcal{R}(B_i)^\ominus. \) Since the first term is in \( A, S_i \in A \) by assumption a).
ii) If \( s_0 < 0 \), \( S_0 = -s_0 B_0 (X_0 - P(X_0|B_0)) - \frac{\epsilon}{K} B_0 \notin A \). 

Suppose by contradiction \( S_0 = -s_0 B_0 (X_0 - (P(X_0|B_0) + \frac{\epsilon}{-s_0 K})) \in A \). 
Since \( \delta B_0 \in R(B_0)^\ast \), \forall \delta > 0 \), assumption a) gives

\[
\frac{S_0}{-s_0} + \delta B_0 = B_0 (X_0 - (P(X_0|B_0) + \frac{\epsilon}{-s_0 K} - \delta)) \in A.
\]

Taking \( 0 < \delta < \frac{\epsilon}{-s_0 K} \), it is \( P(X_0|B_0) + \frac{\epsilon}{-s_0 K} - \delta > P(X_0|B_0) \), so that (25) contradicts the definition of \( P \) in (22).

The gain \( G_2 \) is a function of the gambles \( S_0, S_1 \) and \( K \):

\[
G_2 = \begin{cases} 
S_1 + S_0 - \frac{\epsilon}{K} \sum_{i=0}^{1} B_i & \text{if } s_0 \geq 0 \\
S_1 - S_0 - \frac{\epsilon}{K} \sum_{i \neq 0}^{1} B_i & \text{if } s_0 \leq 0.
\end{cases}
\]

Referring to this representation, define \( T = S_1 + S_0 \) if \( s_0 \geq 0 \), \( T = S_1 - S_0 \), if \( s_0 \leq 0 \). We prove that \( \sup(T|S(\mathfrak{f})) \geq 0 \), distinguishing three cases:

- \( s_1 > 0 \), \( s_0 > 0 \).

By i), \( S_i \in A \), \( i = 0, 1 \). It is also \( S_i \in R(B_i) \), hence \( S_i \in R(B_i) \cap A \), \( i = 0, 1 \). Using assumption c) we deduce \( S_1 + S_0 \in R(B_0 \vee B_1) \setminus R(B_0 \vee B_1)^\ast \), which means that \( \sup(S_1 + S_0|S(\mathfrak{f})) \geq 0 \).

- \( s_0 < 0 \).

By contradiction, suppose \( \sup(S_1 - S_0|S(\mathfrak{f})) = \sup(T|S(\mathfrak{f})) < 0 \), and hence \( \inf(-T|S(\mathfrak{f})) > 0 \). If \( s_1 = 0 \), this means \( -T = S_0 \in R(B_0)^\ast \). If \( s_1 > 0 \), since then \( -T \in R(B_0 \vee B_1) \) by Lemma 2, \( \inf(-T|S(\mathfrak{f})) > 0 \) implies \( -T \in R(B_0 \vee B_1)^\ast \). In both instances, assumption a) can be applied (with \( \lambda = 0 \) if \( s_1 = 0 \), recalling instead that \( S_i \in A \) by i) if \( s_1 > 0 \)) to deduce \( S_0 \in A \), which contradicts ii).

- \( s_1 > 0 \), \( s_0 = 0 \) or \( s_1 = 0 \), \( s_0 > 0 \).

If \( s_1 > 0 \), \( s_0 = 0 \), then using i) \( T = S_1 \in A \cap R(B_1) \). By assumption b), \( \sup(T|S(\mathfrak{f})) = \sup(S_1|B_1) \geq 0 \).

If \( s_1 = 0 \), \( s_0 > 0 \), the argument is analogous.

Recalling (26), \( T = G_2 + \frac{\epsilon}{K} \sum_{i \neq 0}^{1} B_i \). Since \( \sup(T|S(\mathfrak{f})) \geq 0 \) and using the definition of \( K \) at the next equality, we obtain

\[
0 \leq \sup(G_2 + \frac{\epsilon}{K} \sum_{i \neq 0}^{1} B_i|S(\mathfrak{f})) \\
\leq \sup(G_2|S(\mathfrak{f})) + \frac{\epsilon}{K} \sup(\sum_{i \neq 0}^{1} B_i|S(\mathfrak{f})) \\
= \sup(G_2|S(\mathfrak{f}))) + \epsilon.
\]

Thus, \( \sup(G_2|S(\mathfrak{f})) \geq -\epsilon \), \( \forall \epsilon > 0 \), that is \( \sup(G_2|S(\mathfrak{f})) \geq 0 \).

\[\square\]
Unlike the case of coherent conditional lower previsions examined in [22, Section 3.1], \( \mathcal{A} \) does not need to be a cone in Proposition 17: given \( X, Y \in \mathcal{A} \), \( \lambda \geq 0 \), neither \( X + Y \) nor \( \lambda X \) are guaranteed to belong to \( \mathcal{A} \). Actually, condition a) represents a weakening of the cone axioms: if \( X \in \mathcal{A} \), \( Y \in \mathcal{R}(B) \) and \( \lambda \geq 0 \), then \( \lambda X + Y \in \mathcal{A} \). This implies also \( \mathcal{R}(B)^\circ \subseteq \mathcal{A} \forall B \in \mathcal{B} \), a condition that, like also b), is required for coherence as well (see (C1'), (C2') in [22, Section 3.1]).

The interpretation of b) is that of an avoiding partial loss condition: we can expect no gain from owning a gamble in \( \mathcal{R}(B)^\circ \), when \( B \) is true, therefore such gambles cannot be included into \( \mathcal{A} \).

As for c), writing it in an extended form, it tells us that: if \( X_1, X_2 \in \mathcal{A}, B_1 X_1 = X_1, B_2 X_2 = X_2 \), then \((B_1 \lor B_2)(X_1 + X_2) = X_1 + X_2 \) and \( \sup(X_1 + X_2)_{B_1 \lor B_2} \geq 0 \). Note that if \( X_1 \in \mathcal{R}(B_1) \) and \( X_2 \in \mathcal{R}(B_2) \), it always holds that \( X_1 + X_2 \in \mathcal{R}(B_1 \lor B_2) \) by Lemma 2, without having to impose it by means of axiom c).

Therefore, the essential condition in axiom c) is that if \( X_1, X_2 \) are desirable (belonging to \( \mathcal{A} \)), this does not imply that \( X_1 + X_2 \in \mathcal{A} \) (which is required for coherence in [22, 23]), but only that \( X_1 + X_2 \) is not necessarily discarded by resorting to b) with \( B = B_1 \lor B_2 \). To illustrate this concept, let for instance \( B_1 = B_2 = \Omega \) in c), so that \( \mathcal{R}(B_1) = \mathcal{R}(B_2) = \mathcal{R}(B_1 \lor B_2) = \mathcal{R}(\Omega) = \mathcal{X} \). Then, c) implies \( X_1 + X_2 \notin \mathcal{R}(\Omega)^\circ \), making impossible to apply b) in order to discard \( X_1 + X_2 \) from \( \mathcal{A} \).

As for question Q2), an answer is given by the following proposition, when \( \mathcal{P} \) is 2-coherent.

**Proposition 18.** Let \( \mathcal{P} : \mathcal{D}_{LIN} \to \mathbb{R} \) be 2-coherent. Define

\[ \mathcal{A}' = \{ \lambda B(X - x) + Y : X|B \in \mathcal{D}_{LIN}, x < \mathcal{P}(X|B), Y \in \mathcal{X}^\circ, \lambda \geq 0 \} \]

Then the set \( \mathcal{A}' \) is such that:

a') \( a \mathcal{A}' + \mathcal{X}^\circ \subseteq \mathcal{A}' \), \( \forall a \geq 0 \);

b') \( \mathcal{X}^\circ \cap \mathcal{A}' = \{0\} \);

c') \( (\mathcal{A}' + \mathcal{A}') \setminus \{0\} \subseteq \mathcal{X} \setminus \mathcal{X}^\circ \);

d') \( \mathcal{P}(X|B) = \sup\{x \in \mathbb{R} : B(X - x) \in \mathcal{A}'\}, \forall X|B \in \mathcal{D}_{LIN} \).

**Proof.** Proof of a'). Take \( Z_1 = \lambda B(X - x) + Y \in \mathcal{A}', a \geq 0 \) and \( Z_2 \in \mathcal{X}^\circ \). Then, \( aZ_1 + Z_2 = a\lambda B(X - x) + aY + Z_2 \), with \( a\lambda \geq 0, X \in \mathcal{X}, B \in \mathcal{B}^\circ \), \( x \leq \mathcal{P}(X|B) \) and \( aY + Z_2 \in \mathcal{X}^\circ \) (because \( \inf(aY + Z_2) \geq a \inf Y + \inf Z_2 \geq 0 \)). Therefore \( aZ_1 + Z_2 \in \mathcal{A}' \).

Proof of b'). Take \( Z = \lambda B(X - x) + Y \in \mathcal{A}' \).

If \( \lambda = 0 \), \( Z = Y \in \mathcal{X}^\circ \). Therefore \( \sup Z \geq \inf Z \geq 0 \), so that \( Z \notin \mathcal{X}^\circ \) if \( Y \neq 0 \), while \( Z \in \mathcal{X}^\circ \) if and only if \( Y = 0 \).

If \( \lambda > 0 \), \( Z \geq \lambda B(X - x) \) because \( Y \geq \inf Y \geq 0 \). It follows that

\[ \sup(Z|B) \geq \sup(\lambda B(X - x)|B) = \sup(\lambda B(X - \mathcal{P}(X|B)) + \lambda B(\mathcal{P}(X|B) - x)|B) \geq \sup(\lambda B(X - \mathcal{P}(X|B))|B) + \inf(\lambda B(\mathcal{P}(X|B) - x)|B) > 0 \]

27
using the property

$$\sup(X_1 + X_2) \geq \inf X_1 + \sup X_2$$

(27)

with \(X_1 = \lambda B(P(X|B) - x), X_2 = \lambda B(X - P(X|B))\) at the second inequality; the final inequality follows from \(\sup(\lambda B(X - P(X|B))|B) \geq 0\) by 2-coherence of \(P\) (equation (11) with \(s_0 = 0\)) and from \(\inf(\lambda B(P(X|B) - x)|B) > 0\) since \(\lambda > 0, P(X|B) > x\).

The above derivation ensures \(\sup Z \geq \sup Z|B > 0\), i.e. \(Z \not\in \mathcal{X}^\mathbb{Z}\).

Whatever is \(\lambda \geq 0\) then, \(Z \in \mathcal{A}'\) implies either \(Z \not\in \mathcal{X}^\mathbb{Z}\) or \(Z = 0\). Therefore, since \(0 \not\in \mathcal{A}' \cap \mathcal{X}^\mathbb{Z}\), we have that \(\mathcal{A}' \cap \mathcal{X}^\mathbb{Z} = \{0\}\).

Proof of c') To establish c'), we prove that for any \(Z_1, Z_2 \in \mathcal{A}', Z_1 + Z_2 \neq 0\), it holds that \(\sup(Z_1 + Z_2) > 0\).

From the definition of \(\mathcal{A}'\), we have that

\[Z_i = \lambda_i B_i(x_i - x_i) + Y_i \ (i = 1, 2)\]

If \(\lambda_1 = \lambda_2 = 0\), then \(Z_i = Y_i \in \mathcal{X}^\mathbb{Z}, i = 1, 2\), and at least one of \(Y_1, Y_2\) is not zero (because \(Z_1 + Z_2 \neq 0\)). If for instance \(Y_1 \neq 0\), then \(\sup Y_1 > 0\) (because \(Y_1 \in \mathcal{X}^\mathbb{Z}\)). It follows that \(\sup(Z_1 + Z_2) = \sup(Y_1 + Y_2) \geq \sup Y_1 + \inf Y_2 \geq \sup Y_1 > 0\), where the first inequality follows from (27).

We may therefore suppose \(\lambda_1 + \lambda_2 > 0\) in the sequel of the proof, defining \(S(\lambda) = \vee\{B_i : \lambda_i > 0, i = 1, 2\}(\neq \emptyset)\). We have that

\[Z_1 + Z_2|S(\lambda) = \sum_{i=1}^{2} \lambda_i B_i(x_i - x_i)|S(\lambda) + (Y_1 + Y_2)|S(\lambda)\]

\[= \sum_{i=1}^{2} \lambda_i B_i(P(X_i|B_i) - x_i)|S(\lambda) + (Y_1 + Y_2)|S(\lambda)\]

\[\geq \sum_{i=1}^{2} \lambda_i B_i(x_i - P(X_i|B_i))|S(\lambda) + \delta + \inf Y_1 + \inf Y_2,\]

where \(\delta = \min \sum_{i=1}^{2} \lambda_i B_i(P(X_i|B_i) - x_i)|S(\lambda)\). Recalling that \(\lambda_i \geq 0, P(X_i|B_i) > x_i\), \(\lambda_1 + \lambda_2 > 0\), it is easy to realise that \(\delta > 0\).

Using this fact in the strict inequality of the following derivation, we obtain

\[\sup(Z_1 + Z_2) \geq \sup(Z_1 + Z_2|S(\lambda)) \geq\]

\[\sup(\sum_{i=1}^{2} \lambda_i B_i(P(X_i|B_i) - x_i)|S(\lambda)) + \delta + \inf Y_1 + \inf Y_2 >\]

\[\sup(\sum_{i=1}^{2} \lambda_i B_i(x_i - P(X_i|B_i))|S(\lambda)) \geq 0,\]

the final inequality holding because \(P\) is 2-coherent.

Proof of d') Let \(S = \{x : B(X - x) \in \mathcal{A}'\}\).

In the first (and larger) part of the proof, we shall prove that

\[P(X|B) \geq x, \forall x \in S.\]

(28)

For this, let \(\pi \in S\). Therefore,

\[B(X - \pi) = \lambda A(Z - z) + Y,\]

(29)

with \(\lambda \geq 0, A \in \mathcal{B}^\mathbb{Z}, Z \in \mathcal{X}, z < P(Z|A), Y \in \mathcal{X}^\mathbb{Z}\). We distinguish the cases \(\lambda = 0\) and \(\lambda > 0\).
\[ \lambda = 0. \]
From \( \inf(X | B) - \pi = \inf(B(X - \pi) | B) \geq \inf(B(X - \pi)) = \inf Y \geq 0 \) (the last inequality holding because \( Y \in \mathcal{X}^\geq \)), we obtain \( \pi \leq \inf(X | B) \). Therefore also \( \pi \leq \underline{P}(X | B) \), because by 2-coherence \( \inf(X | B) \leq \underline{P}(X | B) \).
(Actually, centered 2-convexity is enough for this, by Proposition 8), a).

\[ \lambda > 0. \]
From (29), \( B(X - \pi) \geq \lambda A(Z - z) \), hence
\[ \sup(\lambda A(Z - z) - B(X - \pi)) \leq 0. \] (30)
Define now
\[
\begin{align*}
X_1 &= \lambda A(z - \underline{P}(Z | A)) - B(\pi - \underline{P}(X | B)) \\
X_2 &= \lambda A(Z - \underline{P}(Z | A)) - B(X - \underline{P}(X | B)) - X_1 \\
&= \lambda A(Z - z) - B(X - \pi).
\end{align*}
\]
Observe that:

i) \( \sup(X_1 + X_2 | A \vee B) \geq 0. \)
Since \( X_1 + X_2 = \lambda A(Z - \underline{P}(Z | A)) - B(X - \underline{P}(X | B)) \), this follows from 2-coherence of \( \underline{P} \) (equation (11), with \( s_1 = \lambda, s_0 = 1 \)).

ii) \( \sup(X_2 | A \vee B) \leq \sup(X_2) \leq 0. \)
In fact, \( X_2 \) is the argument of the supremum in equation (30).
Using i) and ii), we obtain
\[ \sup(X_1 | A \vee B) \geq \sup(X_1 + X_2 | A \vee B) - \sup(X_2 | A \vee B) \geq - \sup(X_2 | A \vee B) \geq 0. \]
Now we know that \( \sup(X_1 | A \vee B) \geq 0 \). On the other hand, \( X_1 | A \vee B \) is a three-valued gamble (at most), and precisely it takes the following values
\[
\begin{align*}
\lambda(z - \underline{P}(Z | A)) - (\pi - \underline{P}(X | B)) &\quad \text{on } A \land B, \text{ when } A \land B \neq \emptyset; \\
-(\pi - \underline{P}(X | B)) &\quad \text{on } \neg A \land B, \text{ when } \neg A \land B \neq \emptyset; \\
\lambda(z - \underline{P}(Z | A)) &< 0 \quad \text{on } A \land \neg B, \text{ when } A \land \neg B \neq \emptyset.
\end{align*}
\]
Therefore, \( -(\pi - \underline{P}(X | B)) \geq \sup(X_1 | A \vee B) \geq 0, \) \( i.e. \) \( \pi \leq \underline{P}(X | B) \).
Thus (28) holds. It remains to observe that \( \forall x < \underline{P}(X | B), \) it is \( B(X - x) \in \mathcal{A} \), by definition of \( \mathcal{A} \) and since \( 0 \in \mathcal{X}^\geq \). This means that \( x \in S \). Consequently
\[ \underline{P}(X | B) = \sup S = \sup \{ x \in \mathbb{R} : B(X - x) \in \mathcal{A} \} \].
\[ \square \]

\footnote{The first inequality can be strict if \( \neg A \land B = \emptyset \). Note that \( \neg A \land B = A \land B = \emptyset \) cannot occur, since it implies \( B = \emptyset \).}
Proposition 18 states the existence of a set of desirable gambles \( \mathcal{A}' \), in accordance with a given 2-coherent conditional lower prevision \( \mathcal{P} \) and satisfying the rationality criteria a'), b'), c'). Comparing a'), b') with a), b) in Proposition 17, a clear similarity comes evident: essentially, the sets \( \mathcal{R}(B)^\succ, \mathcal{R}(B)^\prec, B \in \mathcal{B} \), have been replaced with \( X \succ, X \preceq \) respectively. As a consequence, note that \( 0 \in \mathcal{A}' \).

The interpretation of c') is similar to c) in Proposition 17. It tells that: if \( X_1, X_2 \in \mathcal{A}', X_1 + X_2 \neq 0 \), then \( \sup(X_1 + X_2) > 0 \). Again, coherence would allow the stronger implication \( X_1, X_2 \in \mathcal{A}' \rightarrow X_1 + X_2 \in \mathcal{A}' \), while 2-coherence only ensures that \( X_1 + X_2 \) does not belong to the (near) rejection set \( X \preceq \).

Actually, a'), b'), c') prove to be stronger than a), b), c). This means that any 2-coherent conditional prevision can be represented through a set of desirable gambles \( \mathcal{A}' \) satisfying the necessary axioms a'), b'), c'), but also that, at the same time, \( \mathcal{A}' \) satisfies the weaker axioms a), b), c) in Proposition 17.

6.2 Desirability axioms for 2-convex previsions

A comparison between (4) in Definition 3 and (11) in Definition 5 intuitively suggests that we can get an answer to Q1) for 2-convexity from a reduced form of Proposition 17, with \( \lambda = 1 \). More precisely, the following proposition holds:

**Proposition 19.** Let \( \mathcal{A} \subseteq \mathcal{X} \) be such that
\begin{itemize}
  \item[a)] \( \mathcal{A} + \mathcal{R}(B)^\succ \subseteq \mathcal{A}, \forall B \in \mathcal{B} \);
  \item[b)] \( \mathcal{R}(B)^\prec \cap \mathcal{A} = \emptyset, \forall B \in \mathcal{B} \).
\end{itemize}

Define, \( \forall X|B \in \mathcal{D}_{LIN} \),
\[ P(X|B) = \sup\{x \in \mathbb{R} : B(X - x) \in \mathcal{A}\} \]  
(31)
Then,
\begin{itemize}
  \item[i)] \( P \) is 2-convex on \( \mathcal{D}_{LIN} \);
  \item[ii)] \( P \) is centered iff \( \mathcal{R}(B)^\succ \subseteq \mathcal{A} \forall B \in \mathcal{B} \).
\end{itemize}

**Proof.** Proof of i) The proof is a simplification of that of Proposition 17. Analogously, it is checked that condition (4) in Definition 3 is satisfied for \( \mathcal{G}_{2e} \), where \( \mathcal{P} \) is defined by (31). The same steps are followed: first, the definitions of \( K, S_0, S_1 \) in (23), (24) simplify to
\[ K = \sup(B_0 + B_1|B_0 \lor B_1), \]
\[ S_0 = B_0(X_0 - P(X_0|B_0)) - \frac{\epsilon}{K} B_0, \]
\[ S_1 = B_1(X_1 - P(X_1|B_1)) + \frac{\epsilon}{K} B_1. \]
Then, the following are proven in the same way:
\begin{itemize}
  \item[i)] \( S_1 \in \mathcal{A} \);
  \item[ii)] \( S_0 \notin \mathcal{A} \).
\end{itemize}
Equation (26) reduces here to
\[ G_{2c} = S_1 - S_0 - \epsilon K (B_0 + B_1) \]
and defining \( T = S_1 - S_0 \), it is \( \sup (T | B_0 \lor B_1) \geq 0 \) (see the case \( (s_0 < 0, s_1 > 0) \) in the proof of Proposition 17). This fact is exploited to show that \( \sup (G_{2c} | B_0 \lor B_1) \geq 0 \), with the same computations of the final part in the proof of Proposition 17.

**Proof of 2)** Suppose \( R(B)^\succ \subseteq A \). We prove that then \( P \) is centered. In fact, by (31)
\[ P(0 | B) = \sup \{ x : -Bx \in A \}, \forall 0 | B \in D_{LIN}. \]
For \( x < 0 \), \( \inf (-Bx | B) = -x > 0 \), so that \( -Bx \in R(B)^\succ \subseteq A \).
For \( x > 0 \), \( \sup (-Bx | B) = -x < 0 \), which implies \( -Bx \notin A \) by property b).
Therefore \( \sup \{ x : -Bx \in A \} = 0 \), i.e. \( P(0 | B) = 0 \).
Conversely, suppose now
\[ P(0 | B) = \sup \{ x : -Bx \in A \} = 0, \forall 0 | B \in D_{LIN}. \]
We prove that \( R(B)^\succ \subseteq A \) in two steps.

i) \( -Bx \in A, \forall x < 0 \).
To see this, take \( \bar{x} < 0 \). By definition of supremum, \( \exists \bar{x} : \bar{x} < \bar{x} \leq 0, -B\bar{x} \in A \). Writing \( -B\bar{x} = -B\bar{x} + B(\bar{x} - \bar{x}) \), it is \( B(\bar{x} - \bar{x}) \in R(B)^\succ \), because \( \bar{x} - \bar{x} > 0 \). By property a), \( -Bx \in A + R(B)^\succ \subseteq A \), that is \( -Bx \in A \).

ii) \( R(B)^\succ \subseteq A, \forall B \in B^\sigma \).
For the proof, let \( X \in R(B)^\succ \). This implies \( \inf (X | B) > 0 \), so that \( \delta \) can be chosen, such that \( 0 < \delta < \inf (X | B) \). Then
\[ X = BX = B(X - \inf (X | B) + \delta) - B(\delta - \inf (X | B)). \quad (32) \]
Since \( \delta - \inf (X | B) < 0 \), it is \( -B(\delta - \inf (X | B)) \in A \), by i).
Since \( \inf (B(X - \inf (X | B) + \delta) | B) = \inf (X | B) - \inf (X | B) + \delta > 0 \), it holds that \( B(X - \inf (X | B) + \delta) \in R(B)^\succ \).
Applying axiom a) to the decomposition (32), it ensues that \( X \in A \), that is \( R(B)^\succ \subseteq A \).

\[ \square \]

An analogously reduced form of Proposition 18 allows us to answer question Q2) for 2-convexity.

**Proposition 20.** Let \( \underline{P} : D_{LIN} \to R \) be 2-convex. Define
\[ A' = \{ B(X - x) + Y : X | B \in D_{LIN}, x < P(X | B), Y \in X^\ge \}. \]

31
1) The set $A'$ is such that:
   a) $A' + X^\leq \subseteq A'$;
   b) $A' \cap X^\leq = \emptyset$ if $P$ is 1-AUL;
   c) $P(X|B) = \sup\{x \in \mathbb{R} : B(X - x) \in A'\}$, $\forall X|B \in \mathcal{D}_{LIN}$.

2) If $P$ is centered, then $\mathcal{R}(B)^r \subseteq A' \forall B \in \mathcal{B}$; if $P$ is 1-AUL and $\mathcal{R}(B)^r \subseteq A' \forall B \in \mathcal{B}$, then $P$ is centered.

Proof. Proof of 1). Apart from the converse implication in b), the proof is a simplified version of the proof of Proposition 18. Precisely,

- **Proof of a**). See proof of a') in Proposition 18, with $\lambda = 1$.
- **Proof of b**). If $P$ is 1-AUL, then $A' \cap X^\leq = \emptyset$ follows from the proof of b') in Proposition 18, taking $\lambda = 1$; when proving that $\sup(Z|B) > 0$, the step resorting to 2-coherence uses now 1-AUL to justify by (3) that $\sup(B(X - P(X|B))/|B) \geq 0$.
   We prove now the converse implication, that if $A' \cap X^\leq = \emptyset$ then $P$ is 1-AUL. Suppose $A' \cap X^\leq = \emptyset$ while $P$ is not 1-AUL, which means that there exists $X|B$ such that $\sup(X|B) < P(X|B)$. Then $Z = B(X - \sup(X|B)) \in A'$, because $0 \in X^\leq$. Since $\sup Z = \max(\sup(Z|B), 0)$, it holds that $Z = \max(0, \sup(X|B) - \sup(X|B)) = 0$, it is also $Z \in X^\leq$ and therefore $Z \in A' \cap X^\leq$, contradicting the assumption $A' \cap X^\leq = \emptyset$.
- **Proof of c**). Special case of the proof of d') in Proposition 18 (put $\lambda = 1$ and derive i) from 2-convexity rather than 2-coherence).

Proof of 2). From c), we may write

$$P(0|B) = \sup\{x \in \mathbb{R} : -Bx \in A'\}, \forall 0|B \in \mathcal{D}_{LIN}. \quad (33)$$

We prove that *if $P$ is centered then $\mathcal{R}(B)^r \subseteq A'$.*

Suppose then $P$ centered, which means by (33)

$$P(0|B) = \sup\{x : -Bx \in A'\} = 0, \forall B \in \mathcal{B}.$$  

Let us first prove that

$$-Bx \in A', \forall x < 0. \quad (34)$$

In fact, let $\pi < 0$. By the definition of supremum, $\exists \tilde{x} : \pi < \tilde{x} \leq 0$ and $-B\tilde{x} \in A'$. Hence $-B\tilde{x} = -B\tilde{x} + B(\tilde{x} - \pi) \in A'$ by property a), given that $B(\tilde{x} - \pi) \in X^\leq$, since $\inf(B(\tilde{x} - \pi)) = \min(0, \tilde{x} - \pi) = 0$.

Now let $X \in \mathcal{R}(B)^r$, $\delta : 0 < \delta < \inf(X|B)$. Writing $X = B(X - \inf(X|B) + \delta) - B(\delta - \inf(X|B))$, it holds that $-B(\delta - \inf(X|B)) \in A'$, using (34), and that $B(X - \inf(X|B) + \delta) \in X^\leq$ because $\inf(B(X - \inf(X|B) + \delta)) = \min(0, \inf(X|B) - \inf(X|B) + \delta) = 0$. Therefore $X \in A'$, by property a). Since a generic $X \in \mathcal{R}(B)^r$ has been considered, we have shown that $\mathcal{R}(B)^r \subseteq A'$.
Conversely, let us prove now that if $P$ is $1$-AUL and $\mathcal{R}(B)^\sim \subseteq \mathcal{A}'$, then $P$ is centered.

For this, we show that the supremum in equation (33) is zero, which is equivalent to $P(0|B) = 0$, $\forall 0|B \in \mathcal{D}_{LIN}$.

Suppose $\mathcal{R}(B)^\sim \subseteq \mathcal{A}'$, take $x \in \mathbb{R}$, and consider the gamble $-Bx \in \mathcal{R}(B)$. If $x < 0$ then $\inf(-Bx|B) = -x > 0$, so that $-Bx \in \mathcal{R}(B)^\sim \subseteq \mathcal{A}'$. This implies that the supremum in equation (33) is at least zero.

However, if $x > 0$, it is $\sup(-Bx) \leq 0$, hence $-Bx \in \mathcal{X}^\succeq$. By property b), it follows that $-Bx \notin \mathcal{A}'$ for any positive $x$ and, therefore, the supremum in equation (33) is precisely zero.

Comparing Propositions 17 and 18 with, respectively, Propositions 19 and 20, we note that, in addition to the constraint $\lambda = 1$, 2-convexity requires no condition like c) and c') in Propositions 17 and 18 respectively. Referring, for instance, to c'), this means that, given $X, Y \in \mathcal{A}'$ with $X + Y \neq 0$, 2-convexity does not guarantee $\sup(X + Y) > 0$: summing up two individually desirable gambles could therefore give rise to a partial or even to a sure loss. Moreover, a non-centered 2-convex $P$ suffers from a more serious shortcoming: either it is not even 1-AUL, or $\mathcal{R}(B)^\sim \subseteq \mathcal{A}'$ does not necessarily hold, meaning that a non-negative gamble $X = BX (X \neq 0)$ exists that is considered non-desirable. The main drawbacks of 2-convexity relative to 2-coherence are therefore clearly pointed out also by a comparison through desirability axioms.

7 Conclusions

$N$-convex and $n$-coherent conditional lower previsions broaden the spectrum of uncertainty measures that can be accommodated into a behavioural approach to imprecision, including, for instance, conditional extensions of capacities and niveloids when $n = 2$. This choice for $n$ is the most neatly distinguished from coherence, the other extreme in the spectrum, and that retaining more interesting properties. In particular, centered 2-convex and 2-coherent previsions are stable, meaning that they can be extended on any set preserving their consistency properties. 2-convex and 2-coherent previsions also have a clear meaning in terms of desirability. We believe that the desirability investigation carried out in this paper, although still at a foundational level, is important as it displays first results on how this approach works outside coherence and in the general conditional framework. Further work is necessary to investigate additional properties, like the possible existence of envelope theorems, or properties of already defined notions. In particular, we conjecture that the 2-convex or 2-coherent natural extensions may simplify computing the convex or the coherent natural extensions. As a further generalisation of this work, the consistency notions defined here could be extended to the case of unbounded conditional random variables. This has been done in [20] for coherent conditional lower previsions, while, to the best of our knowledge, a similar investigation for convex conditional previsions is still missing.
Acknowledgements

We wish to thank the referees for their helpful comments.

References


